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A trust region reduced basis Pascoletti-Serafini algorithm for multi-objective PDE-constrained parameter optimization

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Abstract: In the present paper non-convex multi-objective parameter optimization problems are considered which are governed by elliptic parametrized partial differential equations (PDEs). To solve these problems numerically the Pascoletti-Serafini scalarization is applied and the obtained scalar optimization problems are solved by an augmented Lagrangian method. However, due to the PDE constraints, the numerical solution is very expensive so that a model reduction is utilized by using the reduced basis (RB) method. The quality of the RB approximation is ensured by a trust-region strategy which does not require any offline procedure, where the RB functions are computed in a greedy algorithm. Moreover, convergence of the proposed method is guaranteed. Numerical examples illustrate the efficiency of the proposed solution technique.

Keywords: Non-convex multi-objective optimization, partial differential equations, Pascoletti-Serafini method, augmented Lagrangian, reduced basis method, trust-region strategy.

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1. Introduction

Multi-objective optimization plays an important role in many applications, e.g., in industry, medicine or engineering. One of the mentioned examples is the minimization of costs with simultaneous quality optimization in production or the minimization of CO₂ emission in energy generation and simultaneous cost minimization. These problems lead to multiobjective optimization problems (MOPs), where we want to achieve an optimal compromise with respect to all given objectives at the same time. Normally, the different objectives are contradictory such that there exists an infinite number of optimal compromises. The set of these compromises is called the *Pareto set*. The goal is to approximate the Pareto set in an efficient way, which turns out to be more expensive than solving a single objective optimization problem.

Since MOPs are of great importance, there exist several algorithms to solve them. Among the most popular methods are scalarization methods, which transform MOPs into scalar problems. For example, in the weighted sum method [6,19,31], convex combinations of the original objectives are optimized. However, in our case the multi-objective optimization problem

$$\min \hat{J}(u) = (\hat{J}_1(u), \dots, \hat{J}_k(u))^T \quad \text{subject to (s.t.) } u \in \mathcal{U}_{\text{ad}} \quad (\text{MOP})$$

is non-convex with a bounded, non-empty, convex and closed set \mathcal{U}_{ad} . To solve (MOP) a suitable scalarization method in that case is the Pascoletti-Serafini (PS) scalarization [7,22]: For a chosen *reference point* $z \in \mathbb{R}^k$ and a given *target direction* $r \in \mathbb{R}^k$ with $r > 0$ the Pascoletti-Serafini problem is given by

$$\min t \quad \text{s.t. } (t, u) \in \mathbb{R} \times \mathcal{U}_{\text{ad}} \text{ and } \hat{J}(u) - z \leq t r. \quad (\text{P}_{z,r}^{\text{PS}})$$

23 In the present paper ($\mathbf{P}_{z,r}^{\text{PS}}$) is solved by an augmented Lagrangian approach. However,
 24 in our case the evaluation of the objective \hat{J} requires the solution of an elliptic partial
 25 differential equation (PDE) for the given parameter u . This implies further that for
 26 the computation of the gradients $\nabla \hat{J}_i, i = 1, \dots, k$, adjoint PDEs have to be solved; cf.
 27 [13]. Here, surrogate models offer a promising tool to reduce the computational effort
 28 significantly [28]. Examples are dimensional reduction techniques such as the Reduced
 29 Basis (RB) method [12,23]. In an offline phase, a low-dimensional surrogate model of the
 30 PDE is constructed by using, e.g., the greedy algorithm, cf. [2,12,14]. In the online phase,
 31 only the RB model is used to solve the PDE, which saves a lot of computing time.

32 We propose an extension of the method in [1] for solving multi-objective PDE-constrained
 33 parameter optimization problems. This procedure is based on a combination of a trust-
 34 region reduced basis method [3,15] and the PS method. In particular, we discuss different
 35 strategies to handle the increasing number of reduced basis functions, which is crucial
 36 in order to guarantee good performances of the algorithm.

37 The paper is organized as follows: In Section 2 we introduce a general MOP and explain
 38 the PS method, in particular, a hierarchical version of the PS algorithm which turns out
 39 to be very efficient in the numerical realization. The concrete PDE-constrained MOP is
 40 investigated in Section 3. The trust-region RB method and its combination with the PSM
 41 is described in Section 4. Convergence is ensured and the algorithmic realization of the
 42 approach is explained. Numerical examples are discussed in detail in Section 5. Finally,
 43 we draw some conclusions.

44 2. Multi-objective optimization

Let $(\mathcal{U}, \langle \cdot, \cdot \rangle_{\mathcal{U}})$ be a real Hilbert space, $\mathcal{U}_{\text{ad}} \subset \mathcal{U}$ non-empty, convex and closed, $k \geq 2$
 arbitrary and $\hat{J}_1, \dots, \hat{J}_k: \mathcal{U}_{\text{ad}} \subset \mathcal{U} \rightarrow \mathbb{R}$ be given real-valued functions. In this manuscript,
 we assume also that \mathcal{U}_{ad} is bounded. This is an assumption we will require later for the
 convergence of our method. Note that one can derive similar results of this section if \mathcal{U}_{ad}
 is unbounded by introducing additional assumptions; cf. [1]. To shorten the notation,
 we write $\hat{J} := (\hat{J}_1, \dots, \hat{J}_k)^T: \mathcal{U}_{\text{ad}} \rightarrow \mathbb{R}^k$. In the following, we deal with the multi-objective
 optimization problem

$$\min \hat{J}(u) \quad \text{s.t.} \quad u \in \mathcal{U}_{\text{ad}}. \quad (\text{MOP})$$

45 **Definition 1.** a) *The functions $\hat{J}_1, \dots, \hat{J}_k$ are called cost or objective functions. Analo-*
 46 *gously, the vector-valued function $\hat{J}: \mathcal{U}_{\text{ad}} \rightarrow \mathbb{R}^k$ is named the (multi-objective) cost*
 47 *or (multi-objective) objective function.*

48 b) *The Hilbert space \mathcal{U} is named the admissible space, the set \mathcal{U}_{ad} is called the admissible*
 49 *set and a vector $u \in \mathcal{U}_{\text{ad}}$ is called admissible.*

50 c) *The space \mathbb{R}^k is named the objective space and the image set $\hat{J}(\mathcal{U}_{\text{ad}})$ is called the*
 51 *objective set. A vector $y = \hat{J}(u) \in \hat{J}(\mathcal{U}_{\text{ad}})$ is called objective point.*

Definition 2 (Partial ordering on \mathbb{R}^k). *On \mathbb{R}^k we define the partial ordering \leq as*

$$x \leq y : \iff (\forall i \in \{1, \dots, k\}: x_i \leq y_i)$$

for all $x, y \in \mathbb{R}^k$. Moreover, we define

$$x < y : \iff (\forall i \in \{1, \dots, k\}: x_i < y_i).$$

For convenience, we write

$$x \not\leq y : \iff (x \leq y \ \& \ x \neq y)$$

- 52 for all $x, y \in \mathbb{R}^k$ and define the two sets $\mathbb{R}_{\leq}^k := \{y \in \mathbb{R}^k \mid y \leq 0\}$, $\mathbb{R}_{\not\leq}^k := \{y \in \mathbb{R}^k \mid y \not\leq 0\}$.
 53 Analogously, the relations \geq , $>$ and $\not\geq$ as well as the sets \mathbb{R}_{\geq}^k and $\mathbb{R}_{\not\geq}^k$ are defined.

Definition 3 (Pareto optimality). a) An admissible vector $\bar{u} \in \mathcal{U}_{\text{ad}}$ and its corresponding objective point $\bar{y} := \hat{J}(\bar{u}) \in \hat{J}(\mathcal{U}_{\text{ad}})$ are called (locally) weakly Pareto optimal if there is no $\tilde{u} \in \mathcal{U}_{\text{ad}}$ (in a neighborhood of \bar{u}) with $\hat{J}(\tilde{u}) < \hat{J}(\bar{u})$. The sets

$$\mathcal{U}_{\text{opt,w}} := \{u \in \mathcal{U}_{\text{ad}} \mid u \text{ is weakly Pareto optimal}\} \subset \mathcal{U}_{\text{ad}},$$

$$\mathcal{U}_{\text{opt,w,loc}} := \{u \in \mathcal{U}_{\text{ad}} \mid u \text{ is locally weakly Pareto optimal}\} \subset \mathcal{U}_{\text{ad}}$$

are said to be the weak Pareto set and the locally weak Pareto set, respectively. The sets

$$\mathcal{J}_{\text{opt,w}} := \hat{J}(\mathcal{U}_{\text{opt,w}}) \subset \mathbb{R}^k, \quad \mathcal{J}_{\text{opt,w,loc}} := \hat{J}(\mathcal{U}_{\text{opt,w,loc}}) \subset \mathbb{R}^k,$$

- 54 are the weak Pareto front and the locally weak Pareto front, respectively.
 b) An admissible vector $\bar{u} \in \mathcal{U}_{\text{ad}}$ and its corresponding objective point $\bar{y} := \hat{J}(\bar{u}) \in \hat{J}(\mathcal{U}_{\text{ad}})$ are called (locally) Pareto optimal if there is no $\tilde{u} \in \mathcal{U}_{\text{ad}}$ (in a neighborhood of \bar{u}) with $\hat{J}(\tilde{u}) \not\leq \hat{J}(\bar{u})$. The sets

$$\mathcal{U}_{\text{opt}} := \{u \in \mathcal{U}_{\text{ad}} \mid u \text{ is Pareto optimal}\} \subset \mathcal{U}_{\text{ad}},$$

$$\mathcal{U}_{\text{opt,loc}} := \{u \in \mathcal{U}_{\text{ad}} \mid u \text{ is locally Pareto optimal}\} \subset \mathcal{U}_{\text{ad}}$$

are called the Pareto set and the local Pareto set, respectively. The sets

$$\mathcal{J}_{\text{opt}} := \hat{J}(\mathcal{U}_{\text{opt}}) \subset \mathbb{R}^k, \quad \mathcal{J}_{\text{opt,loc}} := \hat{J}(\mathcal{U}_{\text{opt,loc}}) \subset \mathbb{R}^k$$

- 55 are called the Pareto front and the local Pareto front, respectively.

56 If we talk about the different notions of (local) (weak) Pareto optimality in one sentence,
 57 we use the notation $\mathcal{U}_{\text{opt,(w),(loc)}}$ to keep the sentence compact. Analogously, $\mathcal{U}_{\text{opt,(w),loc}}$,
 58 $\mathcal{U}_{\text{opt,(loc)}}$, $\mathcal{J}_{\text{opt,(w),(loc)}}$ etc. are to be understood. An example with the different concept
 59 of Pareto optimality can be found in [1, Example 1.2.6].

60 The next theorem goes back to [5]. It also appears in a similar form in [11,26].

61 **Theorem 1.** Suppose that there is $y \in \hat{J}(\mathcal{U}_{\text{ad}}) + \mathbb{R}_{\geq}^k$ such that the set $(y - \mathbb{R}_{\geq}^k) \cap (\hat{J}(\mathcal{U}_{\text{ad}}) +$
 62 $\mathbb{R}_{\geq}^k)$ is compact. Then it holds $\mathcal{J}_{\text{opt}} \neq \emptyset$.

63 **Proof.** This is a slight generalization of [6, Theorem 2.10] using the argument that adding
 64 \mathbb{R}_{\geq}^k to the set $\hat{J}(\mathcal{U}_{\text{ad}})$ does not change the Pareto front \mathcal{J}_{opt} . \square

65 Given any $y = \hat{J}(u) \in \hat{J}(\mathcal{U}_{\text{ad}})$ with $y \notin \mathcal{J}_{\text{opt}}$, it follows directly from the definition of
 66 Pareto optimality that there is $\bar{y} = \hat{J}(\bar{u}) \in \hat{J}(\mathcal{U}_{\text{ad}})$ with $\bar{y} \not\leq y$. However, even if the
 67 Pareto front \mathcal{J}_{opt} is not empty (e.g., since the assumptions of Theorem 1 are satisfied), it is
 68 not clear that there is $\bar{y} \in \mathcal{J}_{\text{opt}}$ with $\bar{y} \not\leq y$. If this property holds for all $y \in \hat{J}(\mathcal{U}_{\text{ad}}) \setminus \mathcal{J}_{\text{opt}}$,
 69 the set \mathcal{J}_{opt} is said to be *externally stable*; cf. [6,26].

70 **Definition 4.** The set \mathcal{J}_{opt} is said to be externally stable if for every $y \in \hat{J}(\mathcal{U}_{\text{ad}})$ there is
 71 $\bar{y} \in \mathcal{J}_{\text{opt}}$ with $\bar{y} \leq y$. This is equivalent to $\hat{J}(\mathcal{U}_{\text{ad}}) \subset \mathcal{J}_{\text{opt}} + \mathbb{R}_{\geq}^k$.

72 Especially for the investigation of suitable solution methods for solving (MOP), we
 73 are interested in guaranteeing that the Pareto front is externally stable. The next result
 74 provides a sufficient condition for this property.

75 **Theorem 2.** *If for every $y \in \hat{J}(\mathcal{U}_{\text{ad}}) + \mathbb{R}_{\geq}^k$ the set $(y - \mathbb{R}_{\geq}^k) \cap (\hat{J}(\mathcal{U}_{\text{ad}}) + \mathbb{R}_{\geq}^k)$ is compact, then*
 76 \mathcal{J}_{opt} *is externally stable.*

77 **Proof.** For a proof of a similar version of this theorem, we refer to [6, Theorem 2.21]. \square

78 Among the methods to solve multi-objective optimization problems, the ones based on
 79 scalarization techniques are frequently appearing in the literature. Let us mention here
 80 the weighted-sum method [6,31], the Euclidian reference point method [29] and the PS
 81 method [7,22]. Since in our case the set $\hat{J}(\mathcal{U}_{\text{ad}}) + \mathbb{R}_{\geq}^k$ is non-convex, we apply the PS
 82 method which is proven to be able to solve a non-convex (MOP).

83 2.1. The PS method

For a chosen *reference point* $z \in \mathbb{R}^k$ and a given *target direction* $r \in \mathbb{R}_{>}^k$ the PS problem is given by

$$\min t \quad \text{s.t.} \quad (t, u) \in \mathbb{R} \times \mathcal{U}_{\text{ad}} \text{ and } \hat{J}(u) - z \leq t r. \quad (\mathbf{P}_{z,r}^{\text{PS}})$$

Analogously, we can define the PS problem as a scalarization problem. For $z \in \mathbb{R}^k$ and $r \in \mathbb{R}_{>}^k$ we define the scalarization function

$$g_{z,r}: \mathbb{R}^k \rightarrow \mathbb{R}, \quad x \mapsto g_{z,r}(x) := \max_{1 \leq i \leq k} \frac{1}{r_i} (x_i - z_i),$$

and the PS scalarized function

$$\hat{J}^{g_{z,r}}(u) := g_{z,r}(\hat{J}(u)) = \max_{1 \leq i \leq k} \frac{1}{r_i} (\hat{J}_i(u) - z_i) \quad \text{for } u \in \mathcal{U}_{\text{ad}}.$$

Then the reformulated PS problem is given by

$$\min \hat{J}^{g_{z,r}}(u) \quad \text{s.t.} \quad u \in \mathcal{U}_{\text{ad}}. \quad (\mathbf{RP}_{z,r}^{\text{PS}})$$

84 The following theorem proved in [1, Theorem 1.7.3] ensures the equivalence between
 85 $(\mathbf{P}_{z,r}^{\text{PS}})$ and $(\mathbf{RP}_{z,r}^{\text{PS}})$.

86 **Theorem 3.** *Let $z \in \mathbb{R}^k$ and $r \in \mathbb{R}_{>}^k$ be arbitrary. On the one hand, if (\bar{u}, \bar{t}) is a global*
 87 *(local) solution of $(\mathbf{P}_{z,r}^{\text{PS}})$, then \bar{u} is a global (local) solution of $(\mathbf{RP}_{z,r}^{\text{PS}})$ with minimal function*
 88 *value \bar{t} . On the other hand, if \bar{u} is a global (local) solution of $(\mathbf{RP}_{z,r}^{\text{PS}})$, then (\bar{u}, \bar{t}) with $\bar{t} :=$*
 89 $\max_{1 \leq i \leq k} (\hat{J}_i(\bar{u}) - z_i) / r_i$ *is a global (local) solution of $(\mathbf{P}_{z,r}^{\text{PS}})$.*

90 **Assumption 1.** *The cost functions $\hat{J}_1, \dots, \hat{J}_k$ are weakly lower semi-continuous and bounded*
 91 *from below.*

92 **Theorem 4.** *Let Assumption 1 be satisfied and $z \in \mathbb{R}^k$ as well as $r \in \mathbb{R}_{>}^k$ be arbitrary.*
 93 *Then $(\mathbf{RP}_{z,r}^{\text{PS}})$ has a global solution $\bar{u} \in \mathcal{U}_{\text{opt}}$.*

94 **Proof.** A proof of this statement can be found in [1, Corollary 1.7.12]. \square

95 The previous result also shows that the existing global solution of $(\mathbf{RP}_{z,r}^{\text{PS}})$ belongs to the
 96 Pareto set. To guarantee a good reconstruction of the Pareto set by the PS method, one
 97 needs that, given a (weakly) Pareto optimal point, it is possible to choose the parameters

98 z and r such that this point solves $(\mathbf{RP}_{z,r}^{\text{PS}})$. This is stated in [1, Theorem 1.7.13], which we
 99 report here for clearness.

100 **Theorem 5.** *Let $\bar{u} \in \mathcal{U}_{\text{opt},w}$ be arbitrary. Then for every $r \in \mathbb{R}_{>}^k$ and every $\bar{r} \in \mathbb{R}$ we have that*
 101 *\bar{u} is a global solution of $(\mathbf{RP}_{z,r}^{\text{PS}})$ for the reference point $z := \hat{J}(\bar{u}) - \bar{r}$. If even $\bar{u} \in \mathcal{U}_{\text{opt}}$, any*
 102 *other global solution \tilde{u} of $(\mathbf{RP}_{z,r}^{\text{PS}})$ satisfies $\hat{J}(\tilde{u}) = \hat{J}(\bar{u})$.*

103 **Remark 1.** *We refer the reader to [1, Lemma 1.7.15] for the derivation of first-order necessary*
 104 *optimality condition for a global solution of $(\mathbf{P}_{z,r}^{\text{PS}})$.*

105 Thus, the PS method can compute in principle every (locally) (weak) Pareto optimal
 106 point so that many algorithms based on PS method have been proposed. Here we only
 107 mention the ones which are related to (but differ from) our proposed technique. Our
 108 main idea is to keep the parameter r fixed, while varying the reference point z . This
 109 was also proposed in [7], but the method turns out to be not applicable numerically for
 110 $k > 2$. In [21], the authors provide assumptions on the Pareto front to ensure that the
 111 so-called trade-off limits (i.e., points on the Pareto front which cannot be improved in at
 112 least one component), are given by the solution to subproblems. Their idea was then to
 113 find these trade-off points first and then compute the rest of the Pareto front. A similar
 114 idea but with the use of Centroidal Voronoi Tessellation was presented by [20]. Finally,
 115 [16] shows and fixes some problematic behavior associated to the algorithm in [21]. We
 116 follow the idea of the mentioned contributions of hierarchically solving subproblems of
 117 (\mathbf{MOP}) , but with the focus of finding a set of reference points, by looking at subproblems,
 118 for which we can obtain Pareto optimal points. We are then not interested in finding
 119 ‘boundary’ points (i.e., the trade-off limits) of the Pareto front and then filling its ‘interior’
 120 as in [16,20,21], but rather to partly generalize this approach. In what follows, we
 121 characterize which reference points are necessary and/or sufficient for computing the
 122 entire (local) (weak) Pareto front. First we recall the following well-defined mappings;
 123 cf. [1, Definition 1.7.16].

Definition 5. *We define the set-valued mappings*

$$\begin{aligned} \mathcal{Q}_{\text{opt},w} : \mathbb{R}^k &\rightrightarrows \mathcal{U}_{\text{opt},w}, & z &\mapsto \{u \in \mathcal{U}_{\text{ad}} \mid u \text{ is a global solution of } (\mathbf{RP}_{z,r}^{\text{PS}})\}, \\ \mathcal{Q}_{\text{opt},w,\text{loc}} : \mathbb{R}^k &\rightrightarrows \mathcal{U}_{\text{opt},w,\text{loc}}, & z &\mapsto \{u \in \mathcal{U}_{\text{ad}} \mid u \text{ is a local solution of } (\mathbf{RP}_{z,r}^{\text{PS}})\}, \\ \mathcal{Q}_{\text{opt},(\text{loc})} : \mathbb{R}^k &\rightrightarrows \mathcal{U}_{\text{opt},(\text{loc})}, & z &\mapsto \mathcal{Q}_{\text{opt},w,(\text{loc})}(z) \cap \mathcal{U}_{\text{opt},(\text{loc})}. \end{aligned}$$

124 From Theorem 3, it follows that $\mathcal{Q}_{\text{opt},(w),(\text{loc})}(\mathbb{R}^k) = \mathcal{U}_{\text{opt},(w),(\text{loc})}$. Furthermore, if
 125 Assumption 1 is satisfied, we infer from Theorem 4 that $\mathcal{Q}_{\text{opt},(w),(\text{loc})}(z) \neq \emptyset$ for all
 126 $z \in \mathbb{R}^k$. We also introduce the notion of a (locally) (weakly) Pareto sufficient set for the
 127 PSM.

128 **Definition 6.** *A set $Z \subset \mathbb{R}^k$ is called (locally) (weakly) Pareto sufficient if we have*
 129 $\mathcal{Q}_{\text{opt},(w),(\text{loc})}(Z) = \mathcal{U}_{\text{opt},(w),(\text{loc})}$.

130 Hence, a (locally) (weakly) Pareto sufficient set contains the reference points which
 131 allow us to compute the entire (local) (weak) Pareto front. Clearly, the set \mathbb{R}^k is (locally)
 132 (weakly) Pareto sufficient, but this fact is not computationally useful. The next lemma
 133 gives a first condition towards this computational efficiency.

Lemma 1. Let $Z \subset \mathbb{R}^k$ be arbitrary. Z is (locally) (weakly) Pareto sufficient, if

$$\forall \bar{u} \in \mathcal{U}_{\text{opt},(w),(loc)} : \exists t \in \mathbb{R} : \hat{f}(\bar{u}) - tr \in Z. \quad (1)$$

134 **Proof.** Let $Z \subset \mathbb{R}^k$ be such that (1) holds. Let $\bar{u} \in \mathcal{U}_{\text{opt},(w),(loc)}$ be arbitrary. We need to
 135 show that there is a $z \in Z$ with $\bar{u} \in \mathcal{Q}_{\text{opt},(w),(loc)}(z)$. Indeed, by (1) there is $t \in \mathbb{R}$ with
 136 $z := \hat{f}(\bar{u}) - tr \in Z$ and by Theorem 5 we already have $\bar{u} \in \mathcal{Q}_{\text{opt},(w),(loc)}(z)$. \square

137 To proceed we introduce the concepts of ideal point and shifted ideal point, which will
 138 be used to define an optimal Pareto sufficient set¹.

139 **Definition 7.** a) We define the ideal objective point $y^{\text{id}} \in \mathbb{R}^k \cup \{-\infty\}$ by $y_i^{\text{id}} :=$
 140 $\inf_{u \in \mathcal{U}_{\text{ad}}} \hat{f}_i(u)$ for all $i \in \{1, \dots, k\}$.

141 b) For an arbitrary vector $\vec{d} \in \mathbb{R}_{>}^k$ define the shifted ideal point $\vec{y}^{\text{id}} := y^{\text{id}} - \vec{d}$. Let
 142 $D_i \subset \mathbb{R}^k$ be given by $D_i := \{y \in \mathbb{R}^k \mid y \geq \vec{y}^{\text{id}}, y_i = \vec{y}_i^{\text{id}}\}$ for all $i \in \{1, \dots, k\}$. Then
 143 the set $D \subset \mathbb{R}^k$ is defined by $D := \bigcup_{i=1}^k D_i$.

144 c) We define $\mathcal{Z}_{\text{opt},(w),(loc)}^D := \{z \in D \mid \exists \bar{u} \in \mathcal{U}_{\text{opt},(w),(loc)} : \exists t \in \mathbb{R} : z = \hat{f}(\bar{u}) - tr\}$.

145 d) For any $y \in \mathbb{R}^k$ we set $t^D(y) := \min_{i \in \{1, \dots, k\}} (y_i - \vec{y}_i^{\text{id}}) / r_i \in \mathbb{R}$.

Remark 2. It is proved in [1, Lemma 1.7.24] that

$$\mathcal{Z}_{\text{opt},(w),(loc)}^D = \{\hat{f}(\bar{u}) - t^D(\hat{f}(\bar{u}))r \mid \bar{u} \in \mathcal{U}_{\text{opt},(w),(loc)}\}.$$

146 Furthermore, the set $\mathcal{Z}_{\text{opt},(w),(loc)}^D$ is (locally) (weakly) Pareto sufficient and there is a Lipschitz
 147 continuous bijection between $\mathcal{Z}_{\text{opt}}^D$ and the Pareto front \mathcal{J}_{opt} . Unfortunately there is no bijection
 148 between $\mathcal{Z}_{\text{opt},(w),(loc)}^D$ and $\mathcal{J}_{\text{opt},(w),(loc)}$, but the set $\mathcal{Z}_{\text{opt},(w),(loc)}^D$ is still (locally) (weakly) Pareto
 149 sufficient. Therefore, it is anyway possible to use it for the computation of the Pareto front.

150 2.2. Hierarchical PS method

151 Due to Definition 7 and Remark 2 the set $\mathcal{Z}_{\text{opt},(w),(loc)}^D$ can only be computed once the
 152 set $\mathcal{U}_{\text{opt},(w),(loc)}$ is available. Clearly, this characterization of $\mathcal{Z}_{\text{opt},(w),(loc)}^D$ is not useful
 153 for a numerical algorithm. Fortunately, in [1,17] it is shown that the Pareto set has a
 154 hierarchical structure. This means that the (weak) Pareto front and the (weak) Pareto sets
 155 of (MOP) are contained in the set of all (weak) Pareto fronts and (weak) Pareto sets of
 156 all of its subproblems. This particular structure of the Pareto set can be exploited to set
 157 up a hierarchical algorithm for obtaining a superset of $\mathcal{Z}_{\text{opt},(w),(loc)}^D$ without computing
 158 entirely the (local) (weak) Pareto set first.

Definition 8. For the index set $I \subset \{1, \dots, k\}$ we denote by \hat{f}^I the multi-objective cost function
 $(\hat{f}_i)_{i \in I} : \mathcal{U}_{\text{ad}} \rightarrow \mathbb{R}^I$, and call the problem

$$\min \hat{f}^I(u) \quad \text{s.t.} \quad u \in \mathcal{U}_{\text{ad}} \quad (\text{MOP}_I)$$

159 a subproblem of (MOP). For $I, K \subset \{1, \dots, k\}$ with $K \subset I$,

160 a) and for every $y \in \mathbb{R}^I$ we denote by $y^K := (y_i)_{i \in K} \in \mathbb{R}^K$ the canonical projection to \mathbb{R}^K .

161 b) the set $\mathcal{U}_{\text{opt},(w),(loc)}^I := \{u \in \mathcal{U}_{\text{ad}} \mid u \text{ is (loc.) (weak.) Pareto optimal for } (\text{MOP}_I)\}$
 162 denotes the (local) (weak) Pareto set and the set $\mathcal{J}_{\text{opt},(w),(loc)}^I := \hat{f}^I(\mathcal{U}_{\text{opt},(w),(loc)}^I) \subset$
 163 \mathbb{R}^I denotes the (local) (weak) Pareto front of the subproblem (MOP_I).

¹ The word 'optimal' here means that removing any point from the set will cause the loss of the Pareto sufficient property.

c) the (local) (weak) nadir objective point for the subproblem (**MOP_I**) is defined by

$$y_i^{\text{nad},I,(w),(loc)} := \sup\{y_i \mid y \in \mathcal{J}_{\text{opt},(w),(loc)}^I\} \quad \text{for all } i \in I.$$

Definition 9. Let $I \subset \{1, \dots, k\}$ be arbitrary. For a given reference point $z \in \mathbb{R}^{|I|}$ and target direction $r \in \mathbb{R}_{>}^{|I|}$, we define the PS problem for (**MOP_I**) by

$$\min t \quad \text{s.t.} \quad (t, u) \in \mathbb{R} \times \mathcal{U}_{\text{ad}} \text{ and } \hat{f}^I(u) - z \leq tr^I. \quad (\mathbf{P}_{I,z,r}^{\text{PS}})$$

Again, it is possible to show that ($\mathbf{P}_{I,z,r}^{\text{PS}}$) is equivalent (in the sense of Theorem 3) to the problem

$$\min \left(\max_{i \in I} \frac{1}{r_i} (\hat{f}_i(u) - z_i) \right) \quad \text{s.t.} \quad u \in \mathcal{U}_{\text{ad}}. \quad (\mathbf{RP}_{I,z,r}^{\text{PS}})$$

164 Let us mention that the statements proved in Section 2.1 can be adapted for the PS method
165 for the subproblems. Similarly we can also define the sufficient Pareto sets.

Definition 10. Let $I \subset \{1, \dots, k\}$ be arbitrary. Given the vector $\tilde{d} \in \mathbb{R}_{>}^k$ and the shifted ideal point $\tilde{y}^{\text{id}} \in \mathbb{R}^k$, which were both introduced in Definition 7, let $D_i^I \subset \mathbb{R}^I$ be given by

$$D_i^I := \{y \in \mathbb{R}^I \mid y \geq (\tilde{y}^{\text{id}})^I, y_i = \tilde{y}_i^{\text{id}}\} \quad \text{for } i \in I.$$

Then the set $D^I \subset \mathbb{R}^I$ is defined by $D^I := \bigcup_{i \in I} D_i^I$. Moreover, for all $K \subset \{1, \dots, k\}$ we define the sets

$$\mathcal{Z}_{\text{opt},(w),(loc)}^{D^I, K} := \{z \in D^I \mid \exists \bar{u} \in \mathcal{U}_{\text{opt},(w),(loc)}^K : \exists t \in \mathbb{R} : z = \hat{f}^I(\bar{u}) - tr^I\}.$$

166 To ease the notation, we write $\mathcal{Z}_{\text{opt},(w),(loc)}^{D^I} := \mathcal{Z}_{\text{opt},(w),(loc)}^{D^I, I}$. If $I = \{1, \dots, k\}$ we set
167 $\mathcal{Z}_{\text{opt},(w),(loc)}^{D, K} := \mathcal{Z}_{\text{opt},(w),(loc)}^{D^I, K}$ and $\mathcal{Z}_{\text{opt},(w),(loc)}^D := \mathcal{Z}_{\text{opt},(w),(loc)}^{D^I, I}$. Finally, for any $y \in \mathbb{R}^I$
168 we set $t^{D^I}(y) := \min_{i \in I} \frac{y_i - \tilde{y}_i^{\text{id}}}{r_i} \in \mathbb{R}$.

169 Note that Remark 2 can be rewritten for the subproblems. It can be shown that the set
170 $\mathcal{Z}_{\text{opt},(w),(loc)}^{D^I}$ can be computed by using the sets $\mathcal{U}_{\text{opt},(w),(loc)}^K$ for all $K \subsetneq I$. This procedure
171 requires the assumption that the cost functions $\hat{f}_1, \dots, \hat{f}_k$ are upper semi-continuous.
172 Other very technical conditions are omitted to ease and shorten the presentation here.
173 For a reader interested in the details we refer to [1, Sec. 1.7.4.2-1.7.4.4]. Here we just give
174 the necessary numerical condition in order to compute a numerical approximation of
175 the set $\mathcal{Z}_{\text{opt},(w),(loc)}^{D^I}$.

176 To do so, we introduce a grid on D^I as follows

Definition 11. Let $I \subset \{1, \dots, k\}$ be arbitrary. For a given grid size $h > 0$ and any $i \in I$, we define

$$\mathcal{Z}_i^{h,I} := \left\{ z \in D_i^I \mid \forall j \in I \setminus \{i\} : \left(\exists k \geq 0 : z_j = \tilde{y}_j^{\text{id}} + \frac{h}{2} + kh \right) \& \left(z_j \leq y_j^{\text{nad},I,w} - \bar{r}_j^i r_j \right) \right\}.$$

177 Furthermore, we set $\mathcal{Z}^{h,I} := \bigcup_{i \in I} \mathcal{Z}_i^{h,I}$. If $I = \{1, \dots, k\}$, we write $\mathcal{Z}^h := \mathcal{Z}^{h,I}$.

Algorithm 1 Solving (MOP) numerically by the hierarchical PS method

```

1: for  $j = 1 : k$  do
2:   Set  $I := \{j\}$ ;
3:   Compute  $\mathcal{U}_{\text{opt},w}^{\text{num}}(I) = \{u \mid u \text{ minimizes } \hat{f}_j\}$ ;
4:   Choose  $\tilde{d}_j$ , compute  $y_j^{\text{id}}$  and set  $\tilde{y}_j^{\text{id}} = y_j^{\text{id}} - \tilde{d}_j$ ;
5:   Set  $\mathcal{UTZ}^{\text{num}}(I) = \{(u, \tilde{d}_j, \tilde{y}_j^{\text{id}}) \mid u \in \mathcal{U}_{\text{opt},w}^{\text{num}}(I)\}$ ;
6: end for
7: for  $i = 2 : k$  do
8:   for all  $I \subset \{1, \dots, k\}$  with  $|I| = i$  do
9:     Initialize  $\mathcal{U}_{\text{opt},w}^{\text{num}}(I) = \bigcup_{K \subset I} \mathcal{U}_{\text{opt},w}^{\text{num}}(K)$  and  $\mathcal{UTZ}^{\text{num}}(I) = \emptyset$ ;
10:    Compute the reference points  $Z^{\text{num}}(I) = \{z \in \mathcal{Z}^{h,I} \mid \neg(2)\}$ ;
11:    while  $Z^{\text{num}}(I) \neq \emptyset$  do
12:      Choose  $z \in Z^{\text{num}}(I)$  and remove  $z$  from  $Z^{\text{num}}(I)$ ;
13:      Solve  $(\mathbf{P}_{I,z,r}^{\text{PS}})/(\mathbf{RP}_{I,z,r}^{\text{PS}})$ ;
14:      Set  $\mathcal{U}_{\text{opt},w}^{\text{num}}(I) \leftarrow \mathcal{U}_{\text{opt},w}^{\text{num}}(I) \cup \mathcal{Q}_{\text{opt},w}^I(z)$ ;
15:      Set
16:       $\mathcal{UTZ}^{\text{num}}(I) \leftarrow \mathcal{UTZ}^{\text{num}}(I) \cup \{(\bar{u}, \bar{t}, z) \mid (\bar{u}, \bar{t}) \text{ gl. sol. of } (\mathbf{P}_{I,z,r}^{\text{PS}})\}$ ;
17:      Add solutions of PSPs with respect to redundant reference points: Set
18:       $\mathcal{UTZ}^{\text{num}}(I) \leftarrow \mathcal{UTZ}^{\text{num}}(I) \cup \{(\bar{u}, \bar{t}, \bar{z}) \mid (\bar{u}, \bar{t}) \text{ gl. sol. of } (\mathbf{P}_{I,z,r}^{\text{PS}}),$ 
19:       $\bar{z} \in Z^{\text{num}}(I) \cap [z - (\bar{t}r^I - (\hat{f}^I(\bar{u}) - z)), z]\}$ ;
20:      Remove redundant reference points: Set
21:       $Z^{\text{num}}(I) \leftarrow Z^{\text{num}}(I) \setminus [z - (\bar{t}r^I - (\hat{f}^I(\bar{u}) - z)), z]$  for all  $\bar{u} \in \mathcal{Q}_{\text{opt},(w)}^I(z)$ ;
22:    end while
23:   end for
24: end for
25: if computeParetoFront == true then
26:   Remove all  $u \in \mathcal{U}_{\text{opt},w}^{\text{num}}(\{1, \dots, k\})$  with  $u \notin \mathcal{U}_{\text{opt}}$  by a non-dominance test;
27: end if

```

The idea is to only choose reference points that lie on the grid $\mathcal{Z}^{h,I}$ and do not satisfy the condition

$$\exists K \subsetneq I: \exists (\bar{u}, \bar{t}, \bar{z}) \in \mathcal{UTZ}^{\text{num}}(K): z^K = \bar{z}^K \ \& \ z^{I \setminus K} \geq \hat{f}^{I \setminus K}(\bar{u}) - \bar{t}r^{I \setminus K}, \quad (2)$$

178 where $\mathcal{UTZ}^{\text{num}}(K)$ is a numerical approximation of $\mathcal{UTZ}(K) = \{(u, \tilde{d}_j, \tilde{y}_j^{\text{id}}) \mid u \in$
179 $\tilde{\mathcal{U}}_{\text{opt},w}(I)\}$. An explanation for excluding points based on (2) can be found in [1]. Finally,
180 we describe the proposed numerical hierarchical PS method in Algorithm 1.

181 **Remark 3.** In [27], the author introduce three different quality criteria for a scalarization
182 method.

a) **Coverage:** Every part of the Pareto set and front has to be represented in the sets $\mathcal{U}_{\text{opt},w}^{\text{num}}$
and $\mathcal{J}_{\text{opt},w}^{\text{num}}$, respectively. This can be measured by

$$\text{cov}(\mathcal{J}_{\text{opt},(w),(loc)}^{\text{num}}) := \max_{\bar{y} \in \mathcal{J}_{\text{opt},(w),(loc)}^{\text{num}}} \min_{y \in \mathcal{J}_{\text{opt},(w),(loc)}^{\text{num}}} \|\bar{y} - y\|.$$

183 In the case of Algorithm 1, we have that $\text{cov}(\mathcal{J}_{\text{opt},(w),(loc)}^{\text{num}}) = \mathcal{O}(h)$ (cf. [1]).

184 b) **Uniformity:** The points on the Pareto set and front should be distributed (almost) equidis-
185 tantly; cf. [1, Remark 1.7.69-b)].

186 c) **Cardinality:** The number of points contained in the numerical approximation should be
187 reasonable. In the case of Algorithm 1 is not possible to estimate a-priori the number of
188 elements computed by the method. It is possible to show a bound which can be computed
189 when the nadir objective point $y^{\text{nad},(w)}$ is known (cf. [1, Remark 1.7.69-c)]).

190 **3. The non-convex parametric PDE-constrained MOP**

Before defining our exemplary MOP, we introduce the PDE model which will later serve as an equality constraint. Let $\Omega \subset \mathbb{R}^d$, $d \in \{2, 3\}$, be a bounded domain with Lipschitz-continuous boundary $\Gamma = \partial\Omega$. Furthermore, let $\Omega_1, \dots, \Omega_m$ be a pairwise disjoint decomposition of the domain Ω and set $\Gamma_i := \partial\Omega_i \cap \partial\Omega$ for all $i = 1, \dots, m$. Then we are interested in the following elliptic diffusion-reaction equation with Robin boundary condition:

$$-\nabla \cdot \left(\sum_{i=1}^m u_i^k \chi_{\Omega_i}(x) \nabla y(x) \right) + u^r r(x)y(x) = f(x) \quad \text{a.e. in } \Omega, \quad (3a)$$

$$u_i^k \frac{\partial y}{\partial \mathbf{n}}(s) + \alpha y(s) = \alpha y_a(s) \quad \text{a.e. on } \Gamma_i. \quad (3b)$$

For every $i \in \{1, \dots, m\}$, the parameter $u_i^k > 0$ represents the diffusion coefficient on the subdomain Ω_i . By $r \in L^\infty(\Omega)$, we denote a reaction function, which is supposed to satisfy $r > 0$ a.e. in Ω and is controlled by the scalar parameter $u^r > 0$. On the right-hand side of (3a), we have the source term $f \in L^2(\Omega)$. The constant $\alpha > 0$ in (3b) models the heat exchange with the outside of the domain Ω , where a temperature of $y_a \in L^2(\Gamma)$ is assumed. In total, the parameter space is given by $\mathcal{U} = \mathbb{R}^m \times \mathbb{R}$ and any parameter $u \in \mathcal{U}$ can be written as the vector $u = (u^k, u^r)^T$ with $u^k = (u_1^k, \dots, u_m^k)^T \in \mathbb{R}^m$. Setting $H = L^2(\Omega)$ and $V = H^1(\Omega)$ the weak formulation of (3) is

$$a(u; y, \varphi) = \mathcal{F}(\varphi) \quad \text{for all } \varphi \in V \quad (4)$$

for any $u \in \mathcal{U}$. In (4) the parameter-dependent symmetric bilinear form $a(u; \cdot, \cdot) : V \times V \rightarrow \mathbb{R}$ is given by

$$\begin{aligned} a(u; \varphi, \psi) := & \sum_{i=1}^m u_i^k \int_{\Omega_i} \nabla \varphi(x) \cdot \nabla \psi(x) \, dx + u^r \int_{\Omega} r(x) \varphi(x) \psi(x) \, dx \\ & + \alpha \int_{\Gamma} \varphi(s) \psi(s) \, ds \end{aligned}$$

for all $\varphi, \psi \in V$ and $u \in \mathcal{U}$. The linear functional $\mathcal{F} \in V'$ is defined by

$$\mathcal{F}(\varphi) := \int_{\Omega} f(x) \varphi(x) \, dx + \alpha \int_{\Gamma} y_a(s) \varphi(s) \, ds \quad \text{for all } \varphi \in V.$$

Lemma 2. a) For all $u \in \mathcal{U}$ it holds

$$\|a(u; \cdot, \cdot)\|_{L(V, V')} \leq C \|u\|_{\mathcal{U}}$$

191 with a constant $C > 0$, which does not depend on u .

b) For all $u \in \mathcal{U}$ with $u^k > 0$ in \mathbb{R} and $u^r > 0$, it holds

$$a(u; \varphi, \varphi) \geq \min(u_1^k, \dots, u_m^k, u^r) \|\varphi\|_V^2 \quad \text{for all } \varphi \in V.$$

192 c) The mapping $\mathcal{F} \in V'$ is well-defined.

193 **Proof.** All statements follow from similar arguments of [18, Lemma 1.4], where related
194 operators were considered in the parabolic case. \square

Theorem 6. Let $u \in \mathcal{U}$ with $u > 0$ be arbitrary. Then there is a unique solution $y = y(u) \in V$ of (3). Moreover, the estimate

$$\|y\|_V \leq C \left(\|f\|_{L^2(\Omega)} + \|y_a\|_{L^2(\Gamma)} \right) \quad (5)$$

195 holds with a constant $C > 0$, which depends continuously on u , but is independent of f and y_a .

196 **Proof.** The claims follow from the Lax-Milgram theorem (cf. [8]) and Lemma 2. \square

Definition 12. Let $u_{\min}^k \in (0, \infty)^m$ and $u_{\min}^r > 0$ be arbitrary. Then we define the closed set

$$\mathcal{U}_{\text{eq}} := \{u \in \mathcal{U} \mid u^k \geq u_{\min}^k, u^r \geq u_{\min}^r\}.$$

197 In view of Theorem 6, it is possible to define the solution operator $\mathcal{S}: \mathcal{U}_{\text{eq}} \rightarrow V$, which maps
198 any parameter $u \in \mathcal{U}_{\text{eq}}$ to the unique solution $y = \mathcal{S}(u) \in V$ of (4).

199 **Remark 4.** Due to Lemma 2, we can conclude that $a(u; \varphi, \varphi) \geq \alpha_{\min} \|\varphi\|_V^2$ for all $\varphi \in V$ and
200 $u \in \mathcal{U}_{\text{eq}}$, where $\alpha_{\min} := \min((u_{\min}^k)_1, \dots, (u_{\min}^k)_m, u_i^r) > 0$. In particular, the constant C
201 in (5) can be chosen independently of u if we restrict ourselves to parameters $u \in \mathcal{U}_{\text{eq}}$.

Theorem 7. The solution operator $\mathcal{S}: \mathcal{U}_{\text{eq}} \rightarrow V$ is twice continuously Fréchet differentiable. For the first derivative $\mathcal{S}': \mathcal{U}_{\text{eq}} \rightarrow L(\mathcal{U}, V)$, we have that for any $u \in \mathcal{U}_{\text{eq}}$ and $h \in \mathcal{U}$ the function $y^h := \mathcal{S}'(u)h \in V$ solves the equation

$$a(u; y^h, \varphi) = -\partial_u a(u; \mathcal{S}(u), \varphi)h \quad \text{for all } \varphi \in V.$$

The second derivative $\mathcal{S}'': \mathcal{U}_{\text{eq}} \rightarrow L(\mathcal{U}, L(\mathcal{U}, V))$ is given as follows: For any $u \in \mathcal{U}_{\text{eq}}$ and $h_1, h_2 \in \mathcal{U}$, the function $y^{h_1, h_2} := \mathcal{S}''(u)(h_1, h_2)$ solves the equation

$$a(u; y^{h_1, h_2}, \varphi) = -\partial_u a(u; \mathcal{S}'(u)h_1, \varphi)h_2 - \partial_u a(u; \mathcal{S}'(u)h_2, \varphi)h_1 \quad \text{for all } \varphi \in V.$$

Remark 5. By $\partial_u a$ we denote the partial derivative of the mapping a w.r.t. the parameter u . Since a is linear in u , it holds

$$\partial_u a(u; \varphi, \psi)h = a(h; \varphi, \psi), \quad \partial_u^2 a(u; \varphi, \psi) = 0 \in L(\mathcal{U}, \mathcal{U}')$$

for all $u, h \in \mathcal{U}$ and all $\varphi, \psi \in V$. In particular, we can identify $\partial_u a(u; \varphi, \psi) \in \mathcal{U}'$ by

$$\partial_u a(u; \varphi, \psi) = \begin{pmatrix} \int_{\Omega_1} \nabla \varphi(x) \cdot \nabla \psi(x) \, dx \\ \vdots \\ \int_{\Omega_m} \nabla \varphi(x) \cdot \nabla \psi(x) \, dx \\ \int_{\Omega} r(x) \varphi(x) \psi(x) \, dx \end{pmatrix} \in \mathcal{U}$$

202 by using the Riesz representation theorem.

We are now ready to state the multiobjective parametric PDE-constrained optimization problem (MPPOP). Let $k \in \mathbb{N}$ be fixed and

$$\sigma_{\Omega}^{(1)}, \dots, \sigma_{\Omega}^{(k)} \geq 0 \quad \text{as well as} \quad \sigma_{\mathcal{U}}^{(1)}, \dots, \sigma_{\mathcal{U}}^{(k)} \geq 0$$

be non-negative weights. Furthermore, denote by $y_\Omega^{(1)}, \dots, y_\Omega^{(k)} \in H$ the desired states and by $u_d^{(1)}, \dots, u_d^{(k)} \in \mathcal{U}$ the desired parameters. Then we define the multiobjective essential cost functions $\hat{J}_1, \dots, \hat{J}_k: \mathcal{U}_{\text{eq}} \rightarrow \mathbb{R}$ by

$$\hat{J}_i(u) := \frac{\sigma_\Omega^{(i)}}{2} \|\mathcal{S}(u) - y_\Omega^{(i)}\|_H^2 + \frac{\sigma_{\mathcal{U}}^{(i)}}{2} \|u - u_d^{(i)}\|_{\mathcal{U}}^2 \quad \text{for all } u \in \mathcal{U}_{\text{eq}} \text{ and } i \in \{1, \dots, k\}.$$

Moreover, u_a, u_b with $u_a \leq u_b$ are lower and upper bounds on the parameter u which we assume to be finite. We define $\mathcal{U}_{\text{ad}} := \{u \in \mathcal{U} \mid u_a \leq u \leq u_b\}$ and we assume that $\mathcal{U}_{\text{ad}} \subset \mathcal{U}_{\text{eq}}$ holds. Note that \mathcal{U}_{ad} is a closed, convex and bounded set because of the finiteness assumption on u_a and u_b . We are interested in solving

$$\min_{u \in \mathcal{U}_{\text{ad}}} \hat{J}(u) = \min_{u \in \mathcal{U}_{\text{ad}}} (\hat{J}_1(u), \dots, \hat{J}_k(u))^T. \quad (\text{MPPOP})$$

203 Note that, thanks to the assumptions on \mathcal{U}_{ad} and $\sigma_{\mathcal{U}}^{(i)}$, the costs $\hat{J}_1, \dots, \hat{J}_k$ are upper semi-
 204 continuous and Assumption 1 is also satisfied. This problem fits into the framework
 205 of non-convex multiobjective optimization and Algorithm 1 can be applied. The non-
 206 convexity comes from the way the bilinear form depends on the parameter u . This makes,
 207 in fact, the solution mapping non-linear and thus the MPPOP non-convex. To close this
 208 section, we derive the expression of the gradient and Hessian of the cost functionals
 209 $\hat{J}_1, \dots, \hat{J}_k$. We define the i -th adjoint equation and its solution operator as

Definition 13. For $i = 1, \dots, k$, the solution operator of the i -th adjoint equation is $\mathcal{A}_i: \mathcal{U}_{\text{eq}} \rightarrow V$, where for any given $u \in \mathcal{U}_{\text{eq}}$, $p^{(i)} := \mathcal{A}_i(u)$ solves the equation

$$a(u; \varphi, p^{(i)}) = \langle \sigma_\Omega^{(i)} (\mathcal{S}(u) - y_\Omega^{(i)}), \varphi \rangle_H \quad \text{for all } \varphi \in V. \quad (6)$$

210 As shown in [1], this operators satisfy the two following results:

Lemma 3. The solution operator $\mathcal{A}_i: \mathcal{U}_{\text{eq}} \rightarrow V$ is continuously Fréchet differentiable for all $i = 1, \dots, k$. For all $i = 1, \dots, k$, for the first derivative $\mathcal{A}'_i: \mathcal{U}_{\text{eq}} \rightarrow L(\mathcal{U}, V)$, we have that for any $u \in \mathcal{U}_{\text{eq}}$ and $h \in \mathcal{U}$ the function $p^{(i),h} := \mathcal{A}'_i(u)h \in V$ solves the equation

$$a(u; \varphi, p^{(i),h}) = -\partial_u a(u; \varphi, \mathcal{A}_i(u))h + \sigma_\Omega \langle \mathcal{S}'(u)h, \varphi \rangle_{V',V} \quad \text{for all } \varphi \in V. \quad (7)$$

Corollary 1. Let $\mathcal{U}_{\text{ad}} \subset \mathcal{U}_{\text{eq}}$, $u \in \mathcal{U}_{\text{ad}}$ and $h \in \mathcal{U}$ be arbitrary. Then for $i = 1, \dots, k$ the cost functions \hat{J}_i are twice continuously Fréchet differentiable and it holds

$$\begin{aligned} \nabla \hat{J}_i(u) &= -\partial_u a(u; \mathcal{S}(u), \mathcal{A}_i(u)) + \sigma_{\mathcal{U}}(u - u_d^{(i)}) \in \mathcal{U}, \\ \nabla^2 \hat{J}_i(u)h &= -\partial_u a(u; \mathcal{S}'(u)h, \mathcal{A}_i(u)) - \partial_u a(u; \mathcal{S}(u), \mathcal{A}'_i(u)h) + \sigma_{\mathcal{U}}^{(i)}h \in \mathcal{U}. \end{aligned}$$

211 where we use the representation of $\partial_u a(u; \mathcal{S}(u), \mathcal{A}_i(u)) \in \mathcal{U}'$ in \mathcal{U} , cf. Remark 5.

212 3.1. The RB method for MPPOP

One of the limitations of solving the MPPOP directly with the PSM is the high computational cost. Algorithm 1, in fact, requires to solve the state and adjoint equation a large number of times in order to efficiently approximate the Pareto set. Unfortunately, the numerical evaluation of the state and adjoint solution operators is costly due to the high number of degrees of freedom required to apply, for example, the FE method. For this reason, we use the RB method. In the following we explain how the RB method can be applied to our model. From Theorem 6, we know that the weak form of the state

equation admits a unique solution for any control $u \in \mathcal{U}_{\text{eq}}$. This allows us to define the solution operator $\mathcal{S} : \mathcal{U}_{\text{eq}} \rightarrow V$. Now, let us consider the so-called solution manifold $\mathcal{M} := \{\mathcal{S}(u) \mid u \in \mathcal{U}_{\text{eq}}\}$. The goal of the RB method is to provide a low-dimensional subspace $V^\ell \subset V$, which is a good approximation of \mathcal{M} . The subspace V^ℓ is defined as the span of linearly independent snapshots $\mathcal{S}(u_1), \dots, \mathcal{S}(u_\ell)$ for selected parameters $u_1, \dots, u_\ell \in \mathcal{U}_{\text{eq}}$. Clearly, V^ℓ has dimension ℓ and the snapshots constitute its basis. Let us postpone the discussion on how to select good parameters for generating V^ℓ . Given an RB space V^ℓ , we obtain the reduced-order state equation by a Galerkin projection:

$$a(u; y^\ell, \psi) = \mathcal{F}(\psi) \quad \text{for all } \psi \in V^\ell. \quad (8)$$

Also for the reduced-order equation, we have unique solvability for all parameters $u \in \mathcal{U}_{\text{eq}}$. The solution map $\mathcal{S}^\ell : \mathcal{U}_{\text{eq}} \rightarrow V^\ell$, which maps any parameter $u \in \mathcal{U}_{\text{eq}}$ to the unique solution $y^\ell = \mathcal{S}^\ell(u) \in V^\ell$ of (8), is then well-defined. We can similarly define a reduced-order adjoint equation and essential cost functional. For $i = 1, \dots, k$, we define the essential reduced-order cost functions $\hat{J}_i^\ell : \mathcal{U}_{\text{eq}} \rightarrow \mathbb{R}$ by

$$\hat{J}_i^\ell(u) := \frac{\sigma_\Omega^{(i)}}{2} \|\mathcal{S}^\ell(u) - y_\Omega^{(i)}\|_H^2 + \frac{\sigma_u^{(i)}}{2} \|u - u_d^{(i)}\|_U^2,$$

the reduced-order adjoint equation by

$$a(u; \psi, p^{(i),\ell}) = \langle \sigma_\Omega^{(i)} (\mathcal{S}^\ell(u) - y_\Omega^{(i)}), \psi \rangle_H \quad \text{for all } \psi \in V^\ell \quad (9)$$

213 and the reduced-order adjoint solution operator $\mathcal{A}_i^\ell : \mathcal{U}_{\text{eq}} \rightarrow V$. Following Corollary 1,
 214 it is possible to represent the gradient and the Hessian of the essential reduced-order
 215 cost functions \hat{J}_i^ℓ for $i = 1, \dots, k$ by simply replacing the operators \mathcal{S} and \mathcal{A}_i by their
 216 respective reduced-order versions \mathcal{S}^ℓ and \mathcal{A}_i^ℓ . There are still two aspects which remain to
 217 be clarified: first, how to generate an RB space which guarantees a good approximation
 218 of the state and adjoint solution manifolds and, second, how to estimate a-posteriori (i.e.,
 219 without explicitly evaluating the full-order solution operators \mathcal{S} and \mathcal{A}) the error of such
 220 an approximation.

221 For the first aspect, one can think of building an RB space either prior to solving the
 222 reduced-order optimization problem or while solving it. The first approach is the so-
 223 called offline/online decomposition; cf. [9]. This technique exploits a greedy algorithm
 224 in the offline phase, which iteratively searches for the parameter for which the approxi-
 225 mation error between the full- and reduced-order state and adjoint variables is the largest.
 226 Then, the RB space is enriched (by solving the full-order state and adjoint equations at
 227 the respective parameter and orthonormalizing the newly computed snapshots with
 228 respect to the current RB basis) until a pre-defined tolerance for the approximation error
 229 is reached. Once the RB space is computed, the online phase can start: the optimization
 230 problem is solved fast on the reduced-order level. Although this technique is still widely
 231 used in literature, it shows many disadvantages in the context of optimization. At first,
 232 it suffers from the curse of dimensionality: for a high-dimensional parameter space it is
 233 too costly to explore the entire parameter space with a greedy procedure. At second, it is
 234 counter-intuitive to prepare an RB space which is accurate enough for any parameter,
 235 when usually the optimization method follows a (short) pattern in the parameter space
 236 to find the solution or when the Pareto set is contained in some local regions of the
 237 parameter space, as often in the case of non-convex multiobjective problems. Luckily,
 238 the focus has shifted recently towards adapting the RB space while proceeding with
 239 the optimization method. This procedure is followed, e.g., by the methods presented
 240 in [3,15,24,30]. Let us specify that in [3,15,24] the authors proposed and progressively
 241 improved an RB method combined with a TR algorithm, based on more general results
 242 presented in [30]. Such a method constructs the RB space adaptively while the optimizer

243 is computing the optimal solution. Our focus here is on further improving the method
244 in [3], which can be considered the most general among the TR-RB methods.

245 For any of the above-mentioned methods, a-posteriori error estimates are crucial to com-
246 pute upper bounds of the approximation error made by the RB space in reconstructing
247 the solution for a given parameter without any full-order solution at hand. In case of
248 optimization, one is also interested in estimating the error in reconstructing the cost
249 functional and its gradient. For our model, we can use the following estimates:

Theorem 8. *Let $u \in \mathcal{U}_{\text{ad}}$ be arbitrary and denote by $\alpha(u)$ the coercivity constant of the bilinear form $a(u; \cdot, \cdot)$. By Remark 4, it holds $\alpha(u) \geq \alpha_{\min} > 0$. Let the residual $r_{\text{st}}(u; \cdot) \in V'$ be given by $r_{\text{st}}(u; \varphi) := \mathcal{F}(\varphi) - a(u; \mathcal{S}^\ell(u), \varphi)$ for all $\varphi \in V$. Then it holds*

$$\|\mathcal{S}(u) - \mathcal{S}^\ell(u)\|_V \leq \Delta_{\text{st}}(u) := \frac{\|r_{\text{st}}(u; \cdot)\|_{V'}}{\alpha(u)}. \quad (10)$$

For $i = 1, \dots, k$ the residual $r_{\text{adj}}^{(i)}(u; \cdot) \in V'$ of the adjoint equations is given by $r_{\text{adj}}^{(i)}(u; \varphi) := \langle \sigma_\Omega^{(i)}(\mathcal{S}^\ell(u) - y_\Omega^{(i)}), \varphi \rangle_H - a(u; \varphi, \mathcal{A}_i^\ell(u))$ for all $\varphi \in V$. Then it holds

$$\|\mathcal{A}_i(u) - \mathcal{A}_i^\ell(u)\|_V \leq \Delta_{\text{adj}}^{(i)}(u) := \frac{\|r_{\text{adj}}^{(i)}(u; \cdot)\|_{V'} + \sigma_\Omega^{(i)} \Delta_{\text{st}}(u)}{\alpha(u)}.$$

Furthermore, for $i = 1, \dots, k$ we have

$$\begin{aligned} |\hat{J}_i(u) - \hat{J}_i^\ell(u)| &\leq \Delta_{\text{st}}(u) \|r_{\text{adj}}^{(i)}(u; \cdot)\|_{V'} + \sigma_\Omega^{(i)} \Delta_{\text{st}}(u)^2 =: \Delta_{\hat{J}_i^\ell}(u), \\ \|\nabla \hat{J}_i(u) - \nabla \hat{J}_i^\ell(u)\|_{\mathcal{U}} &\leq \|\partial_u a(u; \cdot, \cdot)\| \left(\|\mathcal{S}^\ell(u)\|_V \Delta_{\text{adj}}^{(i)}(u) + \Delta_{\text{st}}(u) \Delta_{\text{adj}}^{(i)}(u) \right. \\ &\quad \left. + \Delta_{\text{st}}(u) \|\mathcal{A}_i^\ell(u)\|_V \right) =: \Delta_{\nabla \hat{J}_i^\ell}(u). \end{aligned}$$

250 **Proof.** A proof of the a-posteriori error estimates for the state and adjoint can be found
251 in [9]. For the cost function and the gradient, we refer to [15, Proposition 2.5]. \square

Note that we only need the reduced-order state and adjoint state to evaluate the a-posteriori error estimates. For our example, the computation of the coercivity constant $\alpha(u)$ is cheap, see Lemma 2. In more general examples, this might not be the case. Thus, one often uses a quickly computable lower bound $\alpha_{\text{LB}}(u)$ instead. Possible methods for computing such a lower bound are, e.g., the min-theta approach (cf. [9]) or the Successive Constraint Method (SCM) (cf. [25]). Note finally that the computation of the terms $\|r_{\text{st}}(u; \cdot)\|_{V'}$ and $\|r_{\text{adj}}^{(i)}(u; \cdot)\|_{V'}$ is not possible in an infinite-dimensional setting. Even after discretization with the FE method, the cost of computing such a term depends on the dimension of the full-order model, which contradicts the request of having a computationally cheap estimate. However, in our case, the parameter-separability of the bilinear form $a(u; \cdot, \cdot)$ can be exploited to preassemble certain quantities in such a way that the computational cost for evaluating $\|r_{\text{st}}(u; \cdot)\|_{V'}$ and $\|r_{\text{adj}}^{(i)}(u; \cdot)\|_{V'}$ only depends on the dimension of the RB space; see, e.g., [25]. Finally, we apply the RB method to (MPPOP): for a given RB space V^ℓ the reduced-order MPPOP reads

$$\min \hat{J}^\ell(u) = (\hat{J}_1^\ell(u), \dots, \hat{J}_k^\ell(u))^T \quad \text{s.t.} \quad u \in \mathcal{U}_{\text{ad}}. \quad (\text{MPPOP}^\ell)$$

For an arbitrary reference point $z \in \mathbb{R}^k$ and target direction $r \in \mathbb{R}^k$, the reduced-order PS problem reads

$$\min_{(u,t)} t \quad \text{s.t.} \quad (t, u) \in \mathbb{R} \times \mathcal{U}_{\text{ad}} \text{ and } \hat{J}_i^\ell(u) - z_i \leq t, \quad i = 1, \dots, k. \quad (\mathbf{P}_{z,r}^{\text{PS},\ell})$$

252 One could then outline an algorithm similar to Algorithm 1 by using an offline/online
 253 splitting. Because of the above-mentioned disadvantages, we focus on combining the
 254 PSP+ with the TR-RB method from [3] and extend it with respect to the method in [1].
 255 The TR method introduces new aspects to the RB implementation, such as the adaptive
 256 construction of the RB space; see next section for further details.

257 4. The TR-RB method

We briefly introduce the method from [3] and clarify how to apply this in combination with the PSM. In Section 4.2 we highlight our extension to this method and how this can reduce the computational time. The basic idea of a TR method is to compute a first-order critical point of a costly optimization problem by iteratively solving some cheap-to-solve approximations in local regions of the admissible space, where these model approximations can be trusted (i.e. are accurate enough). In such a way, one can derive a global method, which converges in a finite number of steps. For each outer iteration $j \geq 0$ of the TR method, the cheap approximation of the objective is generally indicated by $m^{(j)}$ and the trust regions are described by a radius $\delta^{(j)}$. To simplify the exposition, let us stick with the case $\mathcal{U} = \mathbb{R}^m \times \mathbb{R}$, as in Section 3. The TR method solves then, for each $j \geq 0$, the following constrained optimization sub-problems

$$\min_{v \in \mathcal{U}} m^{(j)}(v) \quad \text{s.t.} \quad \|v\|_2 \leq \delta^{(j)}, \tilde{u} := u^{(j)} + v \in \mathcal{U}_{\text{ad}}. \quad (11)$$

Under suitable assumptions, problem (11) admits a unique solution $\bar{v}^{(j)}$, which is used to compute the next outer iteration $u^{(j+1)} = u^{(j)} + \bar{v}^{(j)}$. To further simplify the presentation of the algorithm in [3], let us present it for a general cost functional \mathcal{J} . Later in this section we will give more details about its application to the MPPOP and PSM. The TR-RB version of problem (11) is

$$\min_{\tilde{u} \in \mathcal{U}_{\text{ad}}} \mathcal{J}^{\ell,(j)}(\tilde{u}) \quad \text{s.t.} \quad q^{(j)} := \frac{\Delta_{\mathcal{J}^{\ell,(j)}}(\tilde{u})}{\mathcal{J}^{\ell,(j)}(\tilde{u})} \leq \delta^{(j)}, \quad (12)$$

258 where $\Delta_{\mathcal{J}^{\ell,(j)}}(\tilde{u})$ is an estimate for the error $|\mathcal{J}(\tilde{u}) - \mathcal{J}^{\ell,(j)}(\tilde{u})|$. Looking at (12), one
 259 clearly sees that the role of the model function $m^{(j)}$ is played by the reduced-order model
 260 cost functional. This is perfectly in line with the TR spirit of having a cheap-to-solve
 261 approximation of the original optimization problem. The trust regions are defined
 262 instead through the RB error estimator, which is in fact the way one should use to check
 263 the quality of the approximation. In [15] also the importance of introducing a correction
 264 term on the RB level is discussed to improve the performance of the method. We point
 265 out that this only has to be done if one chooses two separate RB spaces for state and
 266 adjoint equations (see also [3]). This will not be the case for our application. In Algorithm
 267 2, we report the method from [3]. In what follows, we guide the reader through the
 268 features of the algorithm. At first, we need to initialize the reduced-order model at
 269 the initial guess $u^{(0)}$. This means computing $\mathcal{S}(u^{(0)})$ and $\mathcal{A}_i(u^{(0)})$ for $i = 1, \dots, k$ and
 270 generating the RB space $V^{\ell,(0)}$ as their span. Similarly, updating the RB space $V^{\ell,(j)}$ at
 271 the point $u^{(j+1)}$ means computing the full-order quantities $\mathcal{S}(u^{(j+1)})$ and $\mathcal{A}_i(u^{(j+1)})$ for
 272 $i = 1, \dots, k$ and adding them to the RB space by a Gram-Schmidt orthonormalization. In
 273 Line 3 of Algorithm 2, it is required to compute the so-called approximated generalized
 274 Cauchy (AGC) point. We report here its definition according to [15,30].

Algorithm 2 TR-RB algorithm

```

1: Initialize the reduced-order model at  $u^{(0)}$ , set  $j = 0$  and Loop_flag=True;
2: while Loop_flag do
3:   Compute the AGC point  $u_{\text{AGC}}^{(j)}$ ;
4:   Compute  $u^{(j+1)}$  as solution of (12) with stopping criteria (15);
5:   if  $\mathcal{J}^{\ell,(j)}(u^{(j+1)}) + \Delta_{\mathcal{J}^{\ell,(j)}}(u^{(j+1)}) < \mathcal{J}^{\ell,(j)}(u_{\text{AGC}}^{(j)})$  then
6:     Accept  $u^{(j+1)}$ , set  $\delta^{(j+1)} = \delta^{(j)}$ , compute  $\varrho^{(j)}$  and  $g(u^{(j+1)})$ ;
7:     if  $g(u^{(j+1)}) \leq \tau_{\text{FOC}}$  then
8:       Set Loop_flag=False;
9:     else
10:      if  $\varrho^{(j)} \geq \eta_{\varrho}$  then
11:        Enlarge the TR radius  $\delta^{(j+1)} = \beta_1^{-1}\delta^{(j)}$ ;
12:      end if
13:      if not Skip_enrichment_flag( $j$ ) then
14:        Update the RB model at  $u^{(j+1)}$ ;
15:      end if
16:    end if
17:  else if  $\mathcal{J}^{\ell,(j)}(u^{(j+1)}) - \Delta_{\mathcal{J}^{\ell,(j)}}(u^{(j+1)}) > \mathcal{J}^{\ell,(j)}(u_{\text{AGC}}^{(j)})$  then
18:    if  $\beta_1\delta^{(j)} \leq \delta_{\min}$  or Skip_enrichment_flag( $j-1$ ) then
19:      Update the RB model at  $u^{(j+1)}$ ;
20:    end if
21:    Reject  $u^{(j+1)}$ , shrink the radius  $\delta^{(j+1)} = \beta_1\delta^{(j)}$  and go to 4;
22:  else
23:    Compute  $\mathcal{J}(u^{(j+1)})$ ,  $g(u^{(j+1)})$ ,  $\varrho^{(j)}$  and set  $\delta^{(j+1)} = \beta_1^{-1}\delta^{(j)}$ ;
24:    if  $g(u^{(j+1)}) \leq \tau_{\text{FOC}}$  then
25:      Set Loop_flag=False;
26:    else
27:      if Skip_enrichment_flag( $j$ ) and  $\varrho^{(j)} \geq \eta_{\varrho}$  then
28:        Accept  $u^{(j+1)}$ ;
29:      else if  $\mathcal{J}(u^{(j+1)}) \leq \mathcal{J}^{\ell,(j)}(u_{\text{AGC}}^{(j)})$  then
30:        Accept  $u^{(j+1)}$  and update the RB model;
31:        if  $\varrho^{(j)} < \eta_{\varrho}$  then
32:          Set  $\delta^{(j+1)} = \delta^{(j)}$ ;
33:        end if
34:      else
35:        if  $\beta_1\delta^{(j)} \leq \delta_{\min}$  or Skip_enrichment_flag( $j-1$ ) then
36:          Update the RB model at  $u^{(j+1)}$ ;
37:        end if
38:        Reject  $u^{(j+1)}$ , set  $\delta^{(j+1)} = \beta_1\delta^{(j)}$  and go to 4;
39:      end if
40:    end if
41:  end if
42:  Set  $j = j + 1$ ;
43: end while

```

Definition 14. Let $\kappa \in (0, 1)$ and $\kappa_{\text{arm}} \in (0, 1)$ be backtracking parameters. For the current iterate $u^{(j)}$ define $d^{(j)} := \nabla \mathcal{J}^{\ell, (j)}(u^{(j)})$. Let $\alpha \in \mathbb{N}$ be the smallest number for which the two conditions

$$\mathcal{J}^{\ell, (j)}(P_{\mathcal{U}_{\text{ad}}}(u^{(j)} - \kappa^\alpha d^{(j)})) - \mathcal{J}^{\ell, (j)}(u^{(j)}) \leq -\frac{\kappa_{\text{arm}}}{\kappa^\alpha} \|P_{\mathcal{U}_{\text{ad}}}(u^{(j)} - \kappa^\alpha d^{(j)}) - u^{(j)}\|_{\mathcal{U}}^2, \quad (13)$$

$$q^{(j)}(P_{\mathcal{U}_{\text{ad}}}(u^{(j)} - \kappa^\alpha d^{(j)})) \leq \delta^{(j)} \quad (14)$$

275 are satisfied, where $P_{\mathcal{U}_{\text{ad}}}: \mathcal{U} \rightarrow \mathcal{U}_{\text{ad}}$ is the canonical projection onto the closed and convex set
 276 \mathcal{U}_{ad} . Then we define the AGC point as $u_{\text{AGC}}^{(j)} := P_{\mathcal{U}_{\text{ad}}}(u^{(j)} - \kappa^\alpha d^{(j)})$.

The TR-RB subproblem (12) is then solved in Line 4 using a projected Newton-CG algorithm with the AGC point as a warm start and the following termination criteria

$$\|u - P_{\mathcal{U}_{\text{ad}}}(u - \nabla \mathcal{J}^{\ell, (j)}(u))\|_{\mathcal{U}} \leq \tau_{\text{sub}}, \quad \beta_{\text{bound}} \delta^{(j)} \leq q^{(j)}(u) \leq \delta^{(j)}. \quad (15)$$

277 The first condition in (15) is the standard first-order criticality condition with tolerance
 278 $\tau_{\text{sub}} \in (0, 1)$ and the second one was already introduced in [24] to avoid too many
 279 iterations close to the TR boundary, which is generally an area where we are already
 280 starting to trust the model function less. The parameter β_{bound} is usually chosen to be
 281 close to one exactly for this purpose.

An important aspect of TR methods is the decision to accept or reject the step $u^{(j+1)}$. Generally, one asks for the so-called sufficient decrease condition $\mathcal{J}^{\ell, (j+1)}(u^{(j+1)}) \leq \mathcal{J}^{\ell, (j)}(u_{\text{AGC}}^{(j)})$; cf. [30]. Note that this condition requires to update the RB space before being sure that the step will be accepted. If it is rejected, then we performed a costly update without the possibility of exploiting it. Because of this fact, [24] proposed a sufficient (Line 5) and a necessary (Line 17) condition for the sufficient decrease condition. In [15] it is also noted that the full-order quantities $\mathcal{J}(u^{(j+1)})$ and $\nabla \mathcal{J}(u^{(j+1)})$ are cheaply available after updating the RB space. Additionally, [3] introduced the possibility of skipping a redundant enrichment, which is particularly useful at the late stage of the method, where we are close to the optimum. This will prevent the dimension of the RB space from growing too fast, so that the cheap-to-solve property is preserved. The three conditions to be checked in order to decide whether to skip the update of the RB space are contained in the following skipping parameter

$$\begin{aligned} \text{Skip_enrichment_flag}(j) := & (q^{(j)}(u^{(j+1)}) \leq \beta_q \delta^{(j+1)}) \quad \text{and} \\ & \left(\frac{|g(u^{(j+1)}) - g^{\ell, (j)}(u^{(j+1)})|}{g^{\ell, (j)}(u^{(j+1)})} \leq \tau_g \right) \quad \text{and} \\ & \left(\frac{\|\nabla \mathcal{J}^{\ell, (j)}(u^{(j+1)}) - \nabla \mathcal{J}(u^{(j+1)})\|_{\mathcal{U}}}{\|\nabla \mathcal{J}^{\ell, (j)}(u^{(j+1)})\|_{\mathcal{U}}} \leq \min\{\tau_{\text{grad}}, \beta_{\text{grad}} \delta^{(j+1)}\} \right). \end{aligned}$$

where $\beta_q, \beta_{\text{grad}}, \tau_g, \tau_{\text{grad}} \in (0, 1)$ are given parameters and

$$g(u) := \|u - P_{\mathcal{U}_{\text{ad}}}(u - \nabla \mathcal{J}(u))\|_{\mathcal{U}}, \quad g^{\ell, (j)}(u) := \|u - P_{\mathcal{U}_{\text{ad}}}(u - \nabla \mathcal{J}^{\ell, (j)}(u))\|_{\mathcal{U}}.$$

Note also that $g(u) = 0$ is nothing else than the standard first-order condition for optimization problems with constraints on the parameter set. This is the reason why Algorithm 2 terminates when $g(u^{(j+1)}) < \tau_{\text{FOC}}$ holds with $\tau_{\text{FOC}} \in (0, 1)$. For more details on the skipping condition, we refer to [3]. Typically, TR methods also have the option of shrinking (enlarging) the TR radius $\delta^{(j)}$ with some factor $\beta_1 \in (0, 1)$ ($\beta_1^{-1} > 1$,

respectively). In the case of Algorithm 2, we shrink the radius if a point is rejected. We also compute the ratio

$$\rho^{(j)} := \frac{\mathcal{J}(u^{(j)}) - \mathcal{J}(u^{(j+1)})}{\mathcal{J}^{\ell,(j)}(u^{(j)}) - \mathcal{J}^{\ell,(j)}(u^{(j+1)})}.$$

- 282 If this ratio is greater than a parameter $\eta_\rho \in [0.75, 1]$, then the radius is enlarged.
 283 Algorithm 2 is proved to be convergent given some technical assumptions on the problem.
 284 We summarize everything in the following theorem (cf. [3])

Theorem 9. *Suppose that $\mathcal{U}_{\text{ad}} = [u^a, u^b] \subset \mathbb{R}^P$ for some $u^a, u^b \in \mathbb{R}^P$ with $u^a \leq u^b$. Assume that \mathcal{J} and $\mathcal{J}^{\ell,(j)}$ ($j \in \mathbb{N}$) are strictly positive, \mathcal{J} is continuously Fréchet differentiable and $\mathcal{J}^{\ell,(j)}$ is even twice continuously Fréchet differentiable for all $j \in \mathbb{N}$. Moreover, $\nabla \mathcal{J}^{\ell,(j)}$ is uniformly Lipschitz-continuous with respect to j . Suppose that there is $\delta_{\min} > 0$ such that for every $j \in \mathbb{N}$ there exists a TR radius $\delta^{(j)} \geq \delta_{\min}$, for which there is a solution $u^{(j+1)}$ of the TR-RB subproblem (12) which is accepted by Algorithm 2. Assume that the family of functions $(\rho^{(j)})_{j \in \mathbb{N}}$ is uniformly continuous w.r.t. the parameter u and the index j . Then every accumulation point \bar{u} of the sequence of iterates $(u^{(j)})_{j \in \mathbb{N}}$ is a first-order critical point for the full-order optimization problem, i.e., it holds*

$$\|\bar{u} - P_{\mathcal{U}_{\text{ad}}}(\bar{u} - \nabla \mathcal{J}(\bar{u}))\|_{\mathcal{U}} = 0.$$

- 285 In particular, Algorithm 2 terminates after finitely many steps.

- 286 Although many of the assumptions in Theorem 9 are quite technical for the proof, one
 287 can show that they are reasonable in the case of the RB method; cf. [3].

288 4.1. The TR-RB algorithm applied to the PS method

- 289 In this section we show how Algorithm 2 can be applied to the PS method. To this end,
 290 we recall the following lemma from [1].

Lemma 4. *There are constants $C_J, C_{\nabla J}, C_{\nabla^2 J} > 0$ such that for any $j \in \{1, \dots, k\}$, any $u \in \mathcal{U}_{\text{ad}}$ and any choice of the RB space V^ℓ it holds*

$$|\hat{f}_i^\ell(u)| \leq C_J, \quad \|\nabla \hat{f}_i^\ell(u)\|_{\mathcal{U}} \leq C_{\nabla J}, \quad \|\nabla^2 \hat{f}_i^\ell(u)\|_{L(\mathcal{U})} \leq C_{\nabla^2 J}.$$

Lemma 4 immediately implies that the reduced-order gradient is uniformly Lipschitz-continuous with respect to ℓ . We have to solve $(\mathbf{P}_{z,r}^{\text{PS}})$. We follow the approach in [1], where the target direction $r = (1, \dots, 1)$ is chosen and an augmented Lagrangian method is used. Provided a penalty parameter $\mu > 0$, the augmented Lagrangian for $(\mathbf{P}_{z,r}^{\text{PS}})$ is

$$\mathcal{L}_A((u, t, s), \lambda; \mu) := t + \sum_{i=1}^k \lambda_i c_i(u, t, s) + \frac{\mu}{2} \sum_{i=1}^k c_i(u, t, s)^2 \quad (16)$$

with $c_i(u, t, s) = \hat{f}_i(u) - z_i - t + s_i$. The idea is to iteratively solve the subproblems

$$\min \mathcal{L}_A((u, t, s), \lambda; \mu) \quad \text{s.t.} \quad (u, t, s) \in \mathcal{U}_{\text{ad}} \times \mathbb{R} \times \mathbb{R}_{\geq}^k \quad (17)$$

approximately and then update the Lagrange multiplier λ and the penalty parameter μ until the termination criteria

$$\|c(u, t, s)\|_{\mathbb{R}^k} < \tau_{\text{EC}}, \quad (18)$$

$$\|(u, t, s) - P_{\text{ad}}((u, t, s) - \nabla_{(u, t, s)} \mathcal{L}_A((u, t, s), \lambda; \mu))\|_{\mathcal{U} \times \mathbb{R} \times \mathbb{R}^k} < \tau_{\text{FOC}} \quad (19)$$

are satisfied for some tolerances $\tau_{\text{EC}}, \tau_{\text{FOC}} \in (0, 1)$, where $P_{\text{ad}}: \mathcal{U} \times \mathbb{R} \times \mathbb{R}^k \rightarrow \mathcal{U}_{\text{ad}} \times \mathbb{R} \times \mathbb{R}_{\geq}^k$ is the canonical projection onto $\mathcal{U}_{\text{ad}} \times \mathbb{R} \times \mathbb{R}_{\geq}^k$. For further details, we refer to [1, Appendix B]. We want to combine then the augmented Lagrangian method with the TR-RB algorithm to solve problem $(\mathbf{P}_{z,r}^{\text{PS}})$. To do so, we apply Algorithm 2 to solve each subproblem (17). We first define the reduced-order augmented Lagrangian

$$\mathcal{L}_A^\ell((u, t, s), \lambda; \mu) := t + \sum_{i=1}^k \lambda_i c_i^\ell(u, t, s) + \frac{\mu}{2} \sum_{i=1}^k c_i^\ell(u, t, s)^2, \quad (20)$$

with $c_i^\ell(u, t, s) = \hat{f}_i^\ell(u) - z_i - t + s_i$, which leads to the reduced-order subproblem

$$\min \mathcal{L}_A^\ell((u, t, s), \lambda; \mu) \quad \text{s.t.} \quad (u, t, s) \in \mathcal{U}_{\text{ad}} \times \mathbb{R} \times \mathbb{R}_{\geq}^k. \quad (21)$$

Note that in this case the admissible set $\mathcal{U}_{\text{ad}} \times \mathbb{R} \times \mathbb{R}_{\geq}^k$ is unbounded, which collides with the first assumption of Theorem 9. Nevertheless, [1] showed that the $(\mathbf{P}_{z,r}^{\text{PS}})$ problem is also equivalent to

$$\min t \quad \text{s.t.} \quad (t, u) \in [t^{\min}, t^{\max}] \times \mathcal{U}_{\text{ad}} \text{ and } \hat{f}(u) - z \leq t. \quad (22)$$

There is still the problem that the admissible set for the slack variables s is given by $[0, \infty)^k$. However, computing the partial derivative of the augmented Lagrangian \mathcal{L}_A with respect to s_i , we obtain

$$\partial_{s_i} \mathcal{L}_A((u, t, s), \lambda; \mu) = \lambda_i + \mu(\hat{f}_i(u) - z_i - t + s_i) \geq \lambda_i + \mu(-z_i - t^{\max} + s_i).$$

Thus, \mathcal{L}_A is strictly monotonically increasing in s_i for $s_i > -\lambda_i/\mu + z_i + t_{\max} =: s_i^{\max}$. Thus, given the Lagrange multiplier λ and the penalty parameter μ , we can restrict the slack variable s_i to the interval $[0, s_i^{\max}]$. This will not cause any modification to the solvability and the solution of the augmented Lagrangian subproblem. By setting $\mathcal{X}_{\text{ad}} := \mathcal{U}_{\text{ad}} \times [t^{\min}, t^{\max}] \times [0, s^{\max}]$, the equivalent formulation for the augmented Lagrangian subproblem corresponding to (22) reads

$$\min_{(u, t, s) \in \mathcal{X}_{\text{ad}}} \mathcal{L}_A((u, t, s), \lambda; \mu). \quad (23)$$

Similarly, the reduced-order augmented Lagrangian subproblem is given by

$$\min \mathcal{L}_A^\ell((u, t, s), \lambda; \mu) \quad \text{s.t.} \quad (u, t, s) \in \mathcal{X}_{\text{ad}}. \quad (24)$$

Therefore, the goal is to apply Algorithm 2 to solve the subproblem (23). To this end, we define $x = (u, t, s) \in \mathcal{U} \times \mathbb{R} \times \mathbb{R}^k$, $\mathcal{J}(x) = \mathcal{L}_A(x, \lambda; \mu)$ and $\mathcal{J}^{\ell, (j)}(x) = \mathcal{L}_A^{\ell, (j)}(x, \lambda; \mu)$ for any reference point $z \in \mathbb{R}^k$, any Lagrange multiplier $\lambda \in \mathbb{R}_{\geq}^k$ and any penalty parameter

$\mu > 0$. Furthermore, using the a-posteriori estimates of the individual objectives (cf. Theorem 8), we have that

$$\begin{aligned} |\mathcal{J}(x) - \mathcal{J}^{\ell,(j)}(x)| &\leq \sum_{j=1}^k \left(\lambda_j + c |\hat{f}_j^{\ell,(j)}(u) - z_j - t + s_j| \right) \Delta_{\hat{f}_j^{\ell,(j)}}(u) \\ &\quad + \sum_{j=1}^k \frac{c}{2} \left(\Delta_{\hat{f}_j^{\ell,(j)}}(u) \right)^2 =: \Delta_{\mathcal{J}}^{\ell,(j)}(u) \end{aligned}$$

291 for all $u \in \mathcal{U}_{\text{ad}}$, which can be used as a-posteriori error estimate in the TR-RB algorithm.
 292 According to Theorem 9, we still need to show the strict positivity of the costs \mathcal{J} and
 293 $\mathcal{J}^{\ell,(j)}$ and the uniform Lipschitz continuity of the gradient $\nabla \mathcal{J}^{\ell,(j)}$. For the first, we note
 294 that the objectives \mathcal{J} and $\mathcal{J}^{\ell,(j)}$ are bounded from below by $C := t^{\min} - \sum_{i=1}^k \lambda_i^2 / (2\mu_i)$.
 295 Since C depends only on fixed parameters of the optimization problems, we can add
 296 $C + 1$ to the cost functions to obtain strict positivity. Obviously, this will not change the
 297 minimizers. The second property is a bit more technical and we prove it in the following
 298 lemma.

299 **Lemma 5.** *Let the Lagrange multiplier λ and the penalty parameter μ be given. Then the*
 300 *function $\mathcal{J}(\cdot) := \mathcal{L}_A(\cdot, \lambda; \mu)$ is twice continuously Fréchet-differentiable for all $j \in \mathbb{N}$ and the*
 301 *gradient $\nabla \mathcal{J}^{\ell,(j)}$ is uniformly Lipschitz continuous with respect to j .*

Proof. Due to Corollary 1 the cost functions $\hat{f}_1, \dots, \hat{f}_k$ are twice continuously Fréchet-differentiable. Thus, the function $(u, t, s) \mapsto \mathcal{L}_A((u, t, s), \lambda; \mu)$ is also twice continuously Fréchet-differentiable as a composition of twice continuously Fréchet-differentiable functions. Similarly, the reduced-order augmented Lagrangians $\mathcal{L}_A^{\ell,(j)}((\cdot, \cdot, \cdot), \lambda; \mu)$ are also twice continuously Fréchet-differentiable for all $j \in \mathbb{N}$. We have that

$$\begin{aligned} \nabla^2 \mathcal{L}_A^{\ell,(j)}((u, t, s), \lambda; \mu)(h^u, h^t, h^s) = & \\ \left(\begin{array}{c} \sum_{j=1}^k \left((\lambda_j + \mu c_j^{\ell,(j)}) \nabla^2 \hat{f}_j^{\ell,(j)}(u) h^u + \mu (d_j^{\ell,(j)} - h^t + h_j^s) \nabla \hat{f}_j^{\ell,(j)}(u) \right) \\ k\mu h^t - \mu \sum_{j=1}^k (d_j^{\ell,(j)} + h_j^s) \\ \mu (d_1^{\ell,(j)} + h_1^s - h^t) \\ \vdots \\ \mu (d_k^{\ell,(j)} + h_k^s - h^t) \end{array} \right) \end{aligned}$$

302 for any $h = (h^u, h^t, h^s) \in \mathcal{U} \times \mathbb{R} \times \mathbb{R}^k$, where $c_j^{\ell,(j)} := \hat{f}_j^{\ell,(j)}(u) - z_j - t + s_j$ and $d_j^{\ell,(j)} :=$
 303 $\langle \nabla \hat{f}_j^{\ell,(j)}(u), h^u \rangle_{\mathcal{U}}$ for $j \in \{1, \dots, k\}$. Using Lemma 4, we obtain that the Hessian matrix
 304 $\nabla^2 \mathcal{L}_A^{\ell,(j)}((u, t, s), \lambda; \mu)$ can be bounded independently of (u, t, s) and j . Using the mean
 305 value theorem, we can conclude that the gradients $\nabla \mathcal{L}_A^{\ell,(j)}((\cdot, \cdot, \cdot), \lambda; \mu)$ are Lipschitz-
 306 continuous with constant C_L uniformly in j . \square

307 As a consequence of Theorem 9, we have that Algorithm 2 applied to solve the aug-
 308 mented Lagrangian subproblem (23) converges after finitely many steps to a first-order
 309 critical point of (23).

310 **Remark 6.** *Algorithm 2 constructs and updates the RB space during the optimization procedure.*
 311 *In the case of the PS method, we are free to choose what to do for the space constructed during the*
 312 *TR-RB procedure. For example, we can use it for the next augmented Lagrangian subproblem*

313 (and also for the next reference point). We explored different ideas (see also [1]), but we report
314 here only the two most interesting and efficient ones:

- 315 1) Use one common RB space for all the subproblems and reference points, i.e. use a single
316 space V^ℓ for solving the MOP. This strategy acquires efficiency in terms of reconstructing
317 the full-order parameter space during the iteration. Therefore, thanks to the possibility of
318 skipping an enrichment (which is the costly part in Algorithm 2), we expect more and more
319 speed-up, together with accuracy, as the algorithm proceeds.
- 320 2) Use multiple (local) RB spaces. This idea is already exploited by [1,4,10]. In this case, we do
321 not use the previously obtained RB space for the next minimization problem. We generate
322 instead k initial spaces $V_1^\ell, \dots, V_k^\ell$, resulting from the minimization² of the objectives
323 $\hat{J}_1, \dots, \hat{J}_k$. Then at the beginning of every PS problem, we can decide to use the space V_i^ℓ
324 for which $q^{(0)}(u^{(0)}) < \beta_q \delta^{(0)}$ and $\dim V_i^\ell \leq \ell_{\max}$, with $\ell_{\max} \in \mathbb{N}$ being a predefined
325 maximal number of basis functions. If several spaces satisfy these conditions then we select
326 the one for which the value $q^{(0)}(u^{(0)})$ is the smallest. If instead there is no space fulfilling
327 these conditions, we initialize a new space V_{k+1}^ℓ by using the full-order quantities $\mathcal{S}(u^{(0)})$
328 and $\mathcal{A}_i(u^{(0)})$ for $i = 1, \dots, k$.

329 Although these two techniques are already efficient, we noticed that there is a common problem:
330 the number of RB basis functions might grow too fast and prevent a good speed-up for the
331 solution. In particular, this is the case for the first strategy. To fix this issue, we propose different
332 strategies to remove basis functions from V^ℓ in Section 4.2. This approach was not considered in
333 [1,3,15,24] and to our knowledge it has not been addressed in the literature yet. In reduced-order
334 optimization, instead, this is meaningful, since the reduced-order model might grow too fast; see,
335 e.g., [18], in the case of proper orthogonal decomposition.

336 4.2. How to reduce the number of basis functions

337 We point out that what is described in this section can also generally be applied to
338 Algorithm 2 from [3] without any relation to the PS method. Therefore, we use again
339 the general notation \mathcal{J} for the cost, as it was done in the beginning of this section. The
340 methodology to remove a basis function comes from the observation that some basis
341 elements might not be used during the optimization process. Suppose that we start from
342 a point $u^{(0)}$ very far from the optimum. Clearly, after j iterations the point $u^{(j)}$ is in a
343 completely different region of the admissible set compared to the one of the starting
344 point. Hence, the basis functions built for $u^{(0)}$ might give a negligible contribution in
345 spanning the reduced-order model at the point $u^{(j)}$. If this is the case, we can expect that
346 these functions will not play any further role also for the subsequent points and therefore
347 they can be removed to reduce the dimension of the RB space. Our methodologies for
348 removing basis functions are then based on Remark 6 and try to check which basis
349 functions give a negligible contribution for the current iteration of the TR-RB algorithm.
350 Notice that every technique we propose from now on will be applied after updating the
351 RB space in the TR-RB algorithm. The aim is to modify the updated RB space in order to
352 provide a new RB space, where the number of basis functions is reduced.

Technique T1. The first proposed technique is based on the computation of the so-called Fourier coefficients. Given $v \in V$ and a set of orthonormal basis functions $\{\psi_n\}_{n=1}^\ell \subset V^\ell$, the n -th Fourier coefficient is defined as $c_{\mathcal{F}}^{(n)}(v) := \langle v, \psi_n \rangle_V$. Now, T1 consists in computing $c_{\mathcal{F}}^{(n)}(\mathcal{S}(u^{(j+1)}))$ and $c_{\mathcal{F}}^{(n)}(\mathcal{A}_i(u^{(j+1)}))$, $i = 1, \dots, k$, for $n = 1, \dots, \ell$ and remove the basis function ψ_n for which

$$\zeta^{(n)} := \max \left\{ \frac{c_{\mathcal{F}}^{(n)}(\mathcal{S}(u^{(j+1)}))^2}{\sum_{\eta=1}^{\ell} c_{\mathcal{F}}^{(\eta)}(\mathcal{S}(u^{(j+1)}))^2}, \max_{i=1, \dots, k} \left\{ \frac{c_{\mathcal{F}}^{(n)}(\mathcal{A}_i(u^{(j+1)}))^2}{\sum_{\eta=1}^{\ell} c_{\mathcal{F}}^{(\eta)}(\mathcal{A}_i(u^{(j+1)}))^2} \right\} \right\}$$

² Note that this procedure does not require extra computational cost, since we need to solve these problems for the hierarchical PSM anyway

353 is below a certain tolerance. Note, in fact, that the Fourier coefficients indicate the
 354 order of magnitude of the contribution of a given basis function in reconstructing the
 355 new snapshots that we want to add to update the RB. Strategy T1 is also based on
 356 the assumption that the snapshots, which we want to include in an update, are the
 357 most relevant for the new TR subproblem, because they correspond to the last accepted
 358 optimization step $u^{(j+1)}$. The advantage of T1 is that the required Fourier coefficients
 359 are already available from the Gram-Schmidt orthogonalization performed during the
 360 update of the RB space. There is, anyway, a possible drawback of T1 due to the tolerance
 361 we set: it can happen that also important basis functions are removed although one
 362 thinks that the tolerance is small enough. Because of this, we would like to have a criteria
 363 to decide in an unbiased way which basis functions should be removed.

364 **Technique T2.** This approach is based on the idea that once a point $u^{(j+1)}$ is accepted
 365 by the TR-RB algorithm and the RB space is updated, we will compute a provisional
 366 AGC point $u_{\text{AGC}}^{(j+1),\text{prov}}$ (cf. Definition 14) with respect to the previously updated RB space.
 367 One robustness criteria that we demand is that after removing basis functions, this
 368 provisional AGC point is still inside the new TR³, although it might not coincide with
 369 the actual AGC point $u_{\text{AGC}}^{(j+1)}$ that we compute after removing basis functions according
 370 to Line 3 in Algorithm 2⁴. If we do not demand this robustness criteria, we can expect
 371 a deterioration of the TR performances due to lack of accuracy of the RB model in the
 372 steepest descent direction. Another important aspect is to guarantee the convergence of
 373 the TR-RB method, which implies checking that the conditions for accepting the point
 374 $u^{(j+1)}$ are still fulfilled, although we removed basis functions.

In summary, the difference with respect to T1 is then to remove basis functions starting
 from the one with the smallest value of $\zeta^{(n)}$ and proceeding in ascending order until one
 of the following conditions is satisfied

$$\frac{\Delta_{\mathcal{J}^{\ell-\text{rem},(j+1)}}(u_{\text{AGC}}^{(j+1),\text{prov}})}{\mathcal{J}^{\ell-\text{rem},(j+1)}(u_{\text{AGC}}^{(j+1),\text{prov}})} > \beta_q \delta^{(j+1)}, \quad (25a)$$

$$\frac{\Delta_{\nabla \mathcal{J}^{\ell-\text{rem},(j+1)}}(u_{\text{AGC}}^{(j+1),\text{prov}})}{\|\nabla \mathcal{J}^{\ell-\text{rem},(j+1)}(u_{\text{AGC}}^{(j+1),\text{prov}})\|_{\mathcal{U}}} > \min\{\tau_{\text{grad}}, \beta_{\text{grad}} \delta^{(j+1)}\}, \quad (25b)$$

$$\frac{\|\nabla \mathcal{J}^{\ell-\text{rem},(j+1)}(u^{(j+1)}) - \nabla \mathcal{J}(u^{(j+1)})\|_{\mathcal{U}}}{\|\nabla \mathcal{J}^{\ell-\text{rem},(j+1)}(u^{(j+1)})\|_{\mathcal{U}}} > \min\{\tau_{\text{grad}}, \beta_{\text{grad}} \delta^{(j+1)}\}, \quad (25c)$$

$$\frac{|g(u^{(j+1)}) - g^{\ell-\text{rem},(j+1)}(u^{(j+1)})|}{g^{\ell-\text{rem},(j+1)}(u^{(j+1)})} > \tau_g, \quad (25d)$$

$$\mathcal{J}^{\ell-\text{rem},(j+1)}(u^{(j+1)}) > \mathcal{J}^{\ell,(j)}(u_{\text{AGC}}^{(j)}), \quad (25e)$$

$$\mathcal{J}^{\ell-\text{rem},(j+1)}(u_{\text{AGC}}^{(j+1),\text{prov}}) - \mathcal{J}(u^{(j+1)}) > -\kappa_{\text{arm}} \|u_{\text{AGC}}^{(j+1),\text{prov}} - u^{(j+1)}\|_{\mathcal{U}}^2. \quad (25f)$$

375 If one of the conditions (25) holds we re-add the basis function to the RB space and finish
 376 the removal continuing with the TR-RB procedure. T2 is summarized in Algorithm
 377 3.

378 Let us explain the meaning of (25). At first, the superindex $\ell - \text{rem}$ indicates that the
 379 space used to compute the quantity is the RB space obtained after removing a basis
 380 function. Condition (25a) is to check that the provisional AGC point will remain inside
 381 in an accurate-enough region of the TR. Condition (25b) is in the spirit of (25a) but
 382 for the gradient of the objective. Conditions (25c)-(25d) are based on the skipping
 383 enrichment criteria and are checked to ensure convergence and robustness of the method

³ Note that the TR depends on the reduced-order model due to the inequality constraint in (12) and, therefore, changes if we remove basis functions.

⁴ Note that the reduced-order cost function changes by removing a basis function, so that also the first term in (13) differs after this removal.

Algorithm 3 Summary of T2

- 1: Follow the steps in Algorithm 2 until the RB model is updated at $u^{(j+1)}$;
- 2: Compute a provisional AGC point $u_{\text{AGC}}^{(j+1),\text{prov}}$ by using the reduced-order cost function w.r.t. the updated RB model;
- 3: Compute $\zeta^{(n)}$ for $n \in \{1, \dots, \ell\}$;
- 4: **while** None of the conditions in (25) is fulfilled **do**
- 5: Out of all remaining basis functions, remove the one with the smallest value of $\zeta^{(n)}$ from the RB space;
- 6: **end while**
- 7: Add the last removed basis function to the RB space;
- 8: Proceed with Algorithm 2 with the RB space obtained performing Steps 2-7;

384 after the removal. For a similar issue we need to check that the sufficient decrease
 385 condition is fulfilled as well (cf. (25e)). Finally, (25f) is to enforce that the provisional
 386 AGC point is still a Cauchy point. In such a way, we are sure that Algorithm 2 converges
 387 even after performing the basis removal (cf. [3,15]). In this sense, T2 introduces an
 388 unbiased way to deal with the technique introduced in T1. There are still a few aspects
 389 one should comment on before implementing T2. At first, note that all the above-
 390 mentioned conditions are cheaply computable, since they are based either on reduced-
 391 order quantities or the appearing full-order quantities are available because of the
 392 RB update. At second, conditions (25a) and (25b) request efficient and reliable error
 393 estimators. Although for the PSM the efficiency of $\Delta_{\mathcal{J}}^{\ell,(j)}$ is acceptable, it is not the same
 394 for an error estimator $\Delta_{\nabla\mathcal{J}}^{\ell,(j)}$ based on the a-posteriori estimates of the gradients of the
 395 individual objectives. These estimators generally produce a huge overestimation, which
 396 makes them useless in practice. We notice, in fact, that condition (25b) is immediately
 397 triggered in the case of the PSM and we can not remove any basis function. This is the
 398 reason why we solved this issue by two different related approaches:

Technique T2a. We replace the numerator of (25b) by

$$\|\nabla\mathcal{J}^{\ell-\text{rem},(j)}(u_{\text{AGC}}^{(j+1),\text{prov}}) - \nabla\mathcal{J}(u_{\text{AGC}}^{(j+1),\text{prov}})\|_u,$$

399 which is the true error we wanted to estimate, but it is unfortunately costly. It re-
 400 quires the computation of the full-order quantities $\mathcal{S}(u_{\text{AGC}}^{(j+1),\text{prov}})$ and $\mathcal{A}_i(u_{\text{AGC}}^{(j+1),\text{prov}})$,
 401 $i = 1, \dots, k$.

Technique T2b. We replace the numerator of (25b) by

$$\|\nabla\mathcal{J}^{\ell-\text{rem},(j)}(u_{\text{AGC}}^{(j+1),\text{prov}}) - \nabla\mathcal{J}^{\ell,(j+1)}(u_{\text{AGC}}^{(j+1),\text{prov}})\|_u$$

402 which is a cheap approximation of the true error that we suppose to be reliable only after
 403 enough steps of Algorithm 2, however.

404 Clearly, if one has a good estimation of the gradient at hand, T2 can be still used in its
 405 original form.

406 **Technique T3.** Another drawback of T2 is the fact that we first need to remove the
 407 basis function in order to check (25). This implies that when we stop the removal, we
 408 need to add back the last basis function which was removed, because it is containing
 409 important information; cf. Line 7 of Algorithm 3. This results in a waste of time for the
 410 modified Algorithm 2. We decide to add the option of introducing numerical tolerances
 411 for each of the conditions (25). In such a way, the modified algorithm will generally
 412 stop before an important basis function is removed at the price of possibly leaving one
 413 or a few redundant basis functions in the RB space. We think that this is a meaningful

414 modification regarding the time that is wasted reintroducing the removed basis function
 415 into the RB space; cf. Section 5. We indicate this last strategy as T3.

416 5. Numerical experiments

In this section we test Algorithm 2 and compare it with the results obtained in [1, Section 3.2.2]. We use the same numerical setting, which we briefly report here. Let the domain Ω be the two-dimensional unit square, split into four different subdomains $\Omega_1 = (0, 0.5) \times (0, 0.5)$, $\Omega_2 = (0, 0.5) \times (0.5, 1)$, $\Omega_3 = (0.5, 1) \times (0, 0.5)$ and $\Omega_4 = (0.5, 1) \times (0.5, 1)$. For each Ω_i , we consider a corresponding diffusion coefficient $u_i^k \in \mathbb{R}$ in (3) for $i = 1, \dots, 4$. The reaction term $r(x)$ is set to be constantly equal to 1 for any $x \in \Omega$. We impose homogeneous Neumann boundary conditions (i.e., $\alpha = 0$) and a source term $f(x) = \sum_{i=1}^4 c_i \chi_{\Omega_i}(x)$ with $c_1 \approx 2.76$, $c_2 \approx -0.96$, $c_3 \approx 0.51$ and $c_4 \approx -1.66$ generated randomly in order to obtain a problem with a non-convex Pareto front. For the spatial discretization of the state equation, we apply the Finite Element (FE) method with 1340 nodes and piecewise linear basis functions. For (MPPOP) we choose the following three objectives

$$\hat{J}_1(u) := \frac{1}{2} \|\mathcal{S}(u) - y_\Omega^{(1)}\|_H^2 + \frac{\varepsilon}{2} \|u - u_d^{(1)}\|_{\mathcal{U}}^2$$

$$\hat{J}_2(u) := \frac{1}{2} \|\mathcal{S}(u) - y_\Omega^{(2)}\|_H^2 + \frac{\varepsilon}{2} \|u - u_d^{(2)}\|_{\mathcal{U}}^2 \quad \hat{J}_3(u) := \frac{0.05}{2} \|u - u_d^{(3)}\|_{\mathcal{U}}^2$$

with $\varepsilon = 0.002$, the desired states

$$y_\Omega^{(1)}(x) := \chi_{(0,0.5) \times (0,1)}(x), \quad y_\Omega^{(2)}(x) := \chi_{(0.5,1) \times (0,1)}(x),$$

and the desired parameter values

$$u_d^{(1)} = u_d^{(2)} := (2, 0, 0, 0, 0.3)^T, \quad u_d^{(3)} := (2, 1, 1, 1, 0.3)^T.$$

The lower and upper parameter bounds are given by

$$u_a = (2, 0.1, 0.1, 0.1, 0.3)^T \quad \text{and} \quad u_b = (2, 4, 4, 4, 0.3)^T,$$

417 respectively. This implies that $u_1^k = 2$ and $u^r = 0.3$ are seen as constants and we only
 418 optimize over the three parameters u_2^k , u_3^k and u_4^k . Note furthermore, that the desired
 419 parameters $u_d^{(1)} = u_d^{(2)}$ are not admissible. In fact, as for the parameters of the source
 420 term, they were chosen such that the resulting Pareto front is non-convex.

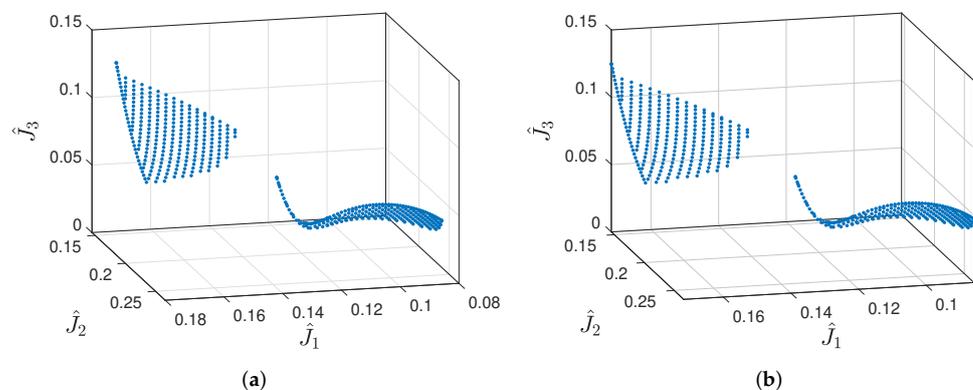


Figure 1. (a) Algorithm 2 no Removal local RB spaces. (b) Algorithm 2 T3 local RB spaces.

421 For the choice of the initial value for PSPs corresponding to reference points for the
 422 entire problem $(\hat{f}_1, \hat{f}_2, \hat{f}_3)$ we do the following: Let \bar{u}^i be the minimizer of \hat{f}_i for $i = 1, 2, 3$.
 423 Recall that the sets D_i have been introduced in Definition 7-(ii). Then, if $z \in D_i$, we
 424 choose \bar{u}^i as the initial value for solving $(P_{z,r}^{PS})$. We additionally choose the shifting
 425 vectors $\vec{d} = 0.001 \cdot (1, 1, 1)^T$, while the grid size h for the reference point grid is set to
 426 $h_{PSM} = 0.003$. For detailed comments and results on the PSM applied on the FE and RB
 427 level, we refer to [1, Section 3.2.2]. We report here only the necessary ones on RB level for
 428 a comparison with our proposed technique. Before doing that, let us mention that the
 429 tolerance chosen in T1 (cf. Section 4.2) for the Fourier coefficient is 10^{-6} . Similarly, we
 430 choose the same tolerance for T3 in order to break the removal algorithm before deleting
 431 important basis functions, i.e. we subtract it on the right-hand side of (25a)-(25f). At first,
 432 to validate our approach, we show in Figure 1 the obtained Pareto fronts by using the
 433 method in [1] (left) and our method (right). As one can see, there is no visible difference.
 434 The approximation error is, in fact, of the order of 10^{-6} for a Pareto point computed by
 435 all the proposed techniques (i.e., T1, T2a, T2b and T3) on average.

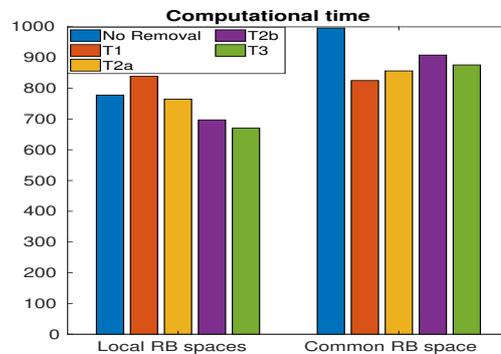


Figure 2. Computational times in seconds for Algorithm 2 with or without basis removal and using the two strategies in Remark 6 for initializing the RB space.

436 In Figure 2 we compare the computational time of Algorithm 2 with all the proposed
 437 techniques to the one of the algorithm in [1]. As one can see, we get a speed-up by using
 438 the proposed techniques in almost all cases. Depending on the strategy from Remark 6,
 439 one technique performs better than the others. Here we try to explain this phenomena
 440 in detail. Let us focus on the common RB space first. In this case, every technique
 441 helps in saving computational time. This is clearly the effect of removing redundant
 442 basis functions, which are particularly frequently included using a large common RB
 443 space. This is the reason why T1 appears to be the most effective, since it is the cheapest
 444 among the techniques (as we said it does not imply additional cost to be checked).
 445 T2a is more robust, but it comes with the price of evaluating the full-order gradient
 446 at the new AGC point and thus results to be slower than T1. Apparently, T2b should
 447 overcome this problem, but the inaccuracy of the RB space in the beginning give a bad
 448 approximation of (25b), resulting in removing too many basis functions which leads to a
 449 worse approximation for the consecutive steps. This worsening of the approximation
 450 results in a way larger number of enrichment steps towards the end of the algorithm,
 451 which also negatively influences the computational time. T3 is comparable with T2a,
 452 meaning that for this example we are removing many basis functions in only a few
 453 instances, rather than frequently removing a few basis functions. Figure 3(b) confirms
 454 the above remarks for the case of a common RB space. In this figure we report the
 455 number of basis functions obtained at the end of Algorithm 2 while this is applied to
 456 compute each Pareto optimal point in the PS method.

457 Now, let us focus on the left group of columns in Figure 2 (and thus on Figure 3(a)),
 458 which corresponds to the computational times in the case of using local RB spaces (cf.

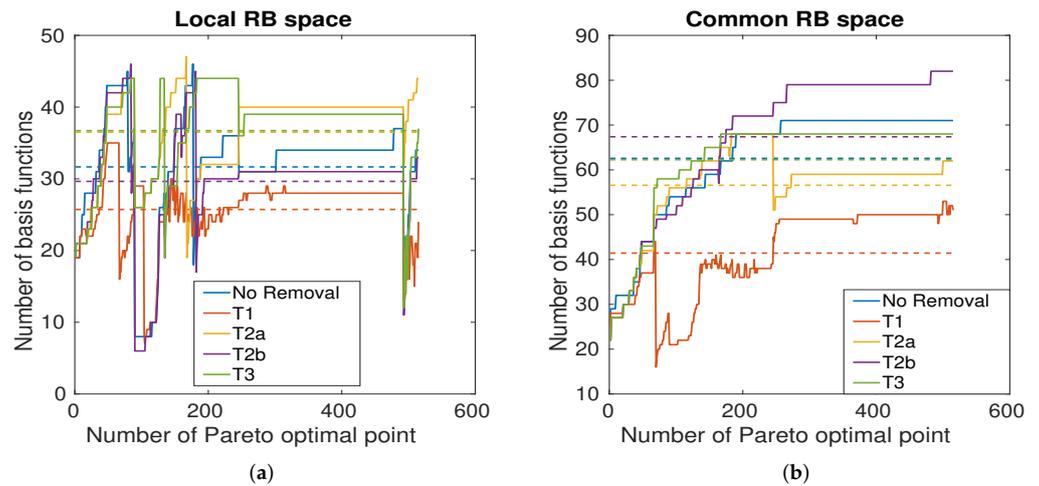


Figure 3. Number of basis functions used to compute each Pareto optimal point. (a) Local RB space. (b) Common RB space. Dashed lines: average number of basis functions.

459 Remark 6). This case is a bit more delicate, since the use of local RB spaces makes it
 460 more difficult to interpret the results. Here the problem of T1 is emerging. The fact
 461 that this technique removes a number of basis functions without any robustness criteria
 462 implies that the method slows down. In the case of local spaces, in fact, we do not have
 463 the same amount of redundant basis functions as it can occur for a common RB space.
 464 Therefore, we should only remove the basis functions which are actually redundant. As
 465 one can note in Figure 3(a), T1 removes a significantly larger amount of basis functions in
 466 comparison to the other techniques. Here the criteria introduced in T2a play their role in
 467 a positive way. We can counteract the effect of T1 in such a way that the computational
 468 time is comparable to the one in [1]. The further simplification introduced in T2b helps
 469 to get an additional speed-up. In contrast to the common RB space, here we have local
 470 spaces which provide a sufficiently good accuracy for approximating (25b) also in the
 471 beginning of the optimization. This is then beneficial for the algorithm, since the cost
 472 of computing the criteria in T2b is way cheaper than T2a, where we need full-order
 473 solves of the state and adjoint equation to compute the gradient at the new AGC point.
 474 Additionally, T3 further improves T2a and T2b in terms of computational time, since in
 475 the case of local RB spaces it is more probable that we indeed remove only a few basis
 476 functions but more frequently than in the case of one common RB space. In this case, it is
 477 important to have tolerances that let us stop before removing an important basis function
 478 and save time for reintroducing it in the RB space. In conclusion, comparing our fastest
 479 method (i.e., Algorithm 2 with local RB spaces and T3) to the slowest (i.e., using [1] with
 480 a common RB space) we get essentially the same results (the approximation error is
 481 10^{-6}) with half of the time, which is roughly 500 seconds. This shows how one should
 482 invest time and resources in providing efficient techniques for reducing the number of
 483 basis functions in the RB space, while using an adaptive TR-RB algorithm. Particularly
 484 in the case of multiobjective optimization, this becomes crucial for a large number of
 485 cost functionals k . To obtain the same resolution of the Pareto front as in Figure 1 for a
 486 large k , we will need to solve the PSPs for many more points, implying higher risk of
 487 having redundant basis functions.

488 6. Conclusions

489 We presented and analyzed novel ways of reducing the dimension of the RB space during
 490 the optimization procedure. To our knowledge, this has not been addressed yet for the
 491 RB method, although it is common for other model order reduction techniques. Such a
 492 removal significantly improved the performances of the TR-RB algorithm in the context
 493 of multiobjective optimization, leading faster to an accurate solution than the already

494 existing techniques. These removal techniques can also be extended to other applications
 495 in which sequential parametric PDE-constrained optimization problems must be solved.
 496 In future work, one can try to achieve further improvements concerning robustness
 497 of the method and deriving tighter a-posteriori error estimators, in particular for the
 498 gradient of the cost function. This is also of great interest in the RB community.

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 511 manuscript, or in the decision to publish the results.

512 Abbreviations

513 The following abbreviations are used in this manuscript:

514

AGC	Approximated generalized Cauchy
CG	Conjugate gradient
FE	Finite element
MOP	Multiobjective optimization problem
MPPOP	Multiobjective parametric PDE-constrained optimization problem
515 PDE	Partial Differential Equation
PS	Pascoletti-Serafini
RB	Reduced basis
s.t.	subject to
TR	Trust-region

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